

Encryption Algorithms & Protocols

Symmetric key Crypto
- **Block cipher**

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Block Cipher (Iterated)

- Plaintext and Ciphertext consist of fixed-sized blocks.
- Ciphertext obtained from plaintext by iterating a round function.
- Input to round function consists of key and output of previous round.
- Usually implemented in software.

Feistel Cipher: Encryption

- Feistel cipher is a type of block cipher, not a specific block cipher.
- Split plaintext block into left and right halves: $P = (L_0, R_0)$
- For each round $i = 1, 2, 3, \dots, n$ compute:
- $L_i = R_{i-1}$
- $R_i = L_{i-1} \oplus F(R_{i-1}, K_i)$, where F is round function and K_i is **subkey**
- Ciphertext: $C = (L_n, R_n)$.

Feistel Cipher: Decryption

- Start with Ciphertext: $C = (L_n, R_n)$.
- For each round $i = n, n - 1, \dots, 2, 1$ compute:
- $R_{i-1} = L_i$
- $L_{i-1} = R_i \oplus F(R_{i-1}, K_i)$, where F is round function and K_i is **subkey**
- Plaintext: $P = (L_0, R_0)$.
- Formula “works” for any function F
- But only secure for certain functions F

Feistel Cipher: Example

- Plain text [0 1 1 1 1 0 1 0 0 0 0 1].
- key = [1st to 2nd -- 2nd to 3rd -- 3rd to 1st]
- $L_0 = [0\ 1\ 1\ 1\ 1\ 0]$, $R_0 = [\underline{1\ 0\ 0}\ \underline{0\ 0\ 1}]$,
- **1st iteration :**
- $L_1 = R_0 = [1\ 0\ 0\ 0\ 0\ 1]$, $R_0 = [1\ 0\ 0\ 0\ 0\ 1]$,
- $F(R_0, K_1) = F(R_0, K_1) = [\underline{0\ 1\ 0}\ \underline{1\ 0\ 0}]$ [1st to 2nd -- 2nd to 3rd -- 3rd to 1st]
- $R_1 = L_0 \oplus F(R_0, K_1) = [0\ 1\ 1\ 1\ 1\ 0] \oplus [0\ 1\ 0\ 1\ 0\ 0] = [0\ 0\ 1\ 0\ 1\ 0]$
- **Ciphertext = [1 0 0 0 0 1 0 0 1 0 1 0]**

**Same procedure
can go to the 2nd,
3rd and so on
iterations**

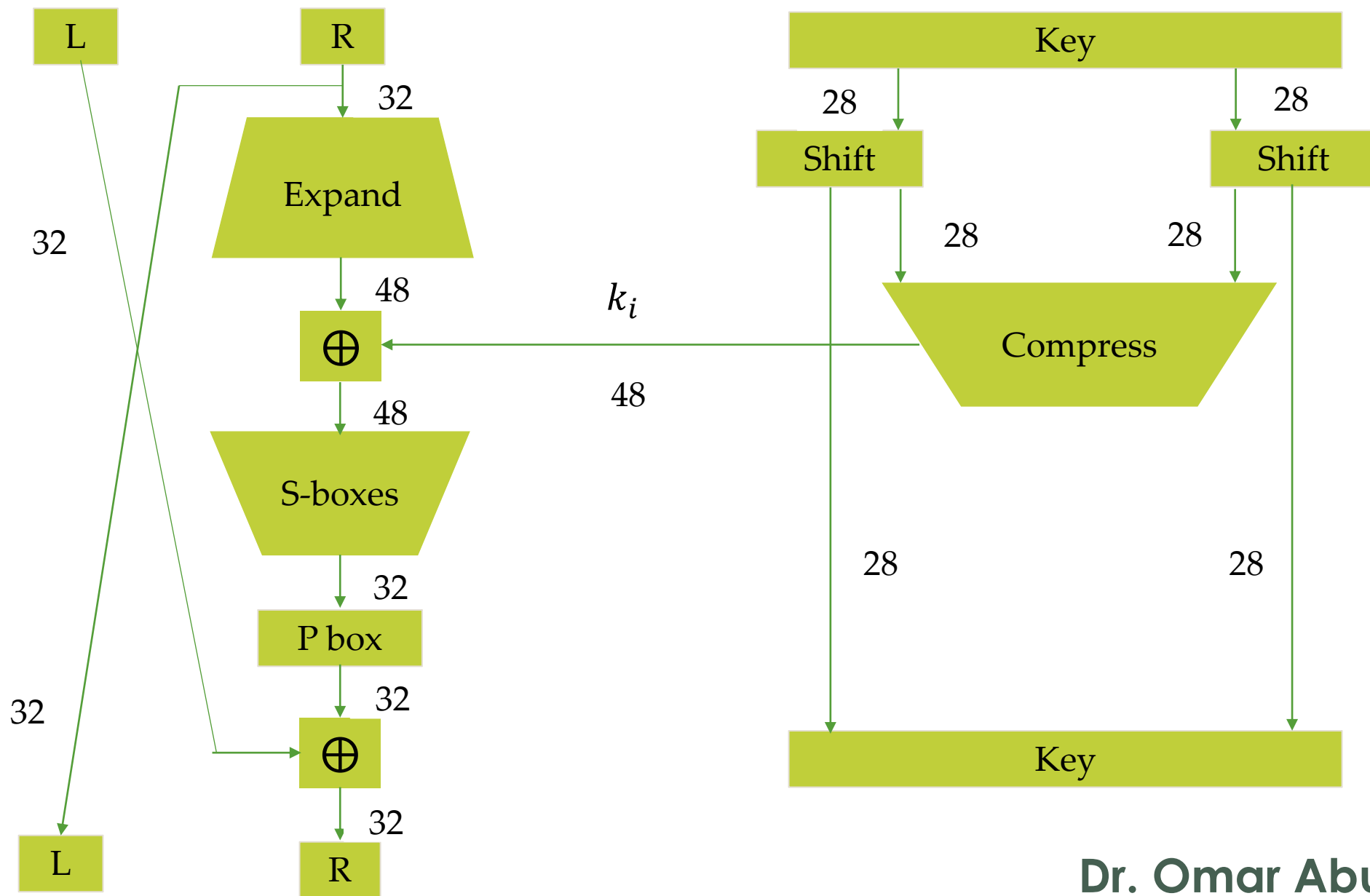
Data Encryption Standard

- DES developed in 1970's.
- Based on IBM's Lucifer cipher.
- DES was U.S. government standard.
- DES development was controversial.
 - NSA secretly involved.
 - Design process was secret.
 - Key length reduced from 128 to 56 bits.
 - Subtle changes to Lucifer algorithm.

DES Numerology

- DES is a Feistel cipher with...
 - 64 bit block length.
 - 56 bit key length.
 - 16 rounds.
 - 48 bits of key used each round (Subkey).
- Each round is simple (for a block cipher).
- Security depends heavily on “S-boxes”.
 - Each S-boxes maps 6 bits to 4 bits.

One Round of DES



Triple DES

- Today, 56 bit DES key is too small.
 - Exhaustive key search is feasible.
- But DES is everywhere, so what to do?
- Triple DES or 3DES (112 bit key).
 - $C = E(D(E(P, K_1), K_2), K_1)$.
 - $P = D(E(D(C, K_1), K_2), K_1)$.
- Encrypt-Decrypt-Encrypt with 2 keys

Advanced Encryption Standard AES

- The AES algorithm, also known as (Rijndael Algorithm) is a symmetric block cipher algorithm that takes block size of 128 bits and converts them into ciphertext using Key of 128, 192 or 256 bits (independent of block size)
- AES performs on byte data, instead of bit data.
- Number of rounds depends on Key length
 - 128 bit Key length uses 10 rounds
 - 192 bit Key length uses 12 rounds
 - 256 bit Key length uses 14 rounds

DES vs AES

DES

- Key length 56 bits
- Block size 64 bits
- Fixed number of rounds (16)
- Implemented slower

AES

- Key length 128/192/256 bits
- Block size 128 bits
- Number of rounds dependent on key length
- Implemented faster

A Few Other Block Ciphers

- **International Data Encryption Algorithm (IDEA)**
- Invented by James Massey
- One of the giants of modern crypto
- IDEA has 64-bit block, 128-bit key
- IDEA uses mixed-mode arithmetic
- Combine different math operations
 - IDEA the first to use this approach
 - Frequently used today

Blowfish

- Blowfish encrypts 64-bit blocks
- Key is variable length, up to 448 bits
- Invented by Bruce Schneier
- Almost a Feistel cipher
- $R_i = L_{i-1} \oplus K_i$
- $L_i = R_{i-1} \oplus F(L_{i-1} \oplus K_i)$
- The round function F uses 4 S-boxes
 - Each S-box maps 8 bits to 32 bits
- Key-dependent S-boxes
 - S-boxes determined by the key

... Thank you ...

