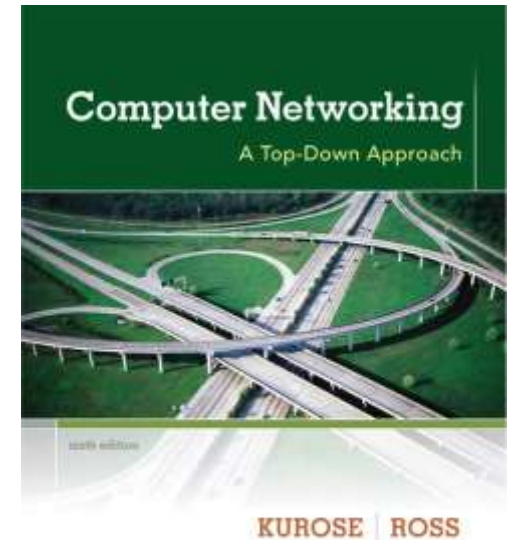


Chapter 1 Introduction

Lecture 2



*Computer
Networking: A Top
Down Approach*
6th edition
Jim Kurose, Keith Ross
Addison-Wesley
March 2012

2021_2020

Chapter 1: roadmap

1.1 what is the Internet?

1.2 network edge

- end systems, access networks, links

1.3 network core

- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

1.5 protocol layers, service models

1.6 networks under attack: security

1.7 history

A closer look at network structure:

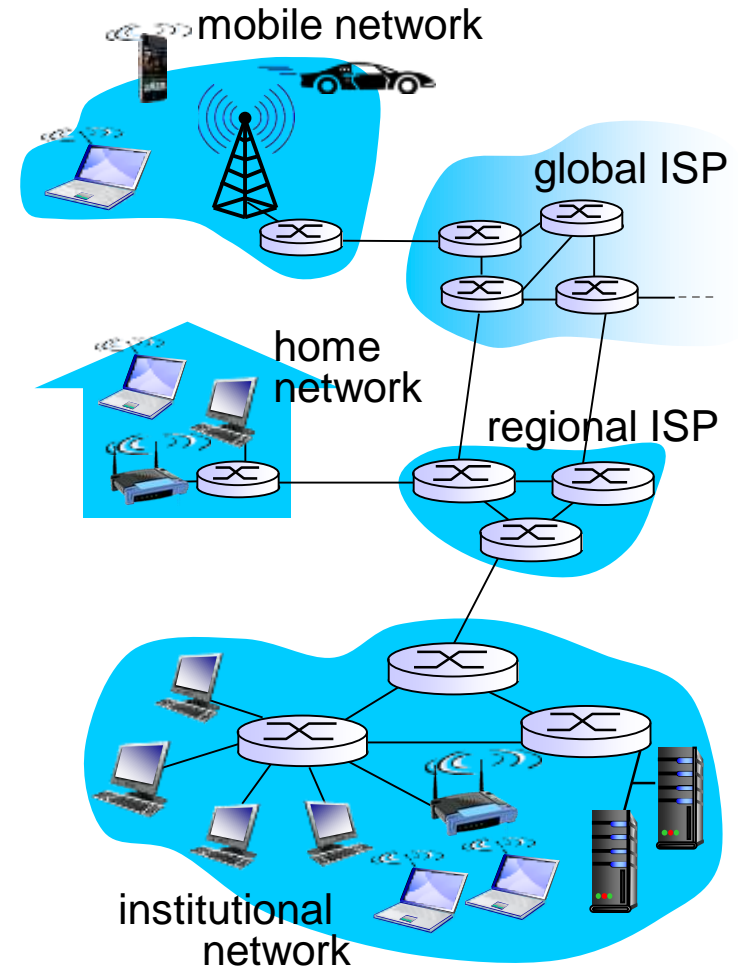
❖ *network edge:*

- hosts: clients and servers
- servers often in data centers

❖ *access networks, physical media:* wired, wireless communication links

❖ *network core:*

- interconnected routers
- network of networks



end systems:

❖ **End systems** are also referred to as *hosts* because they host (that is, run) *application* programs such as a Web browser program, a Web server program, an e-mail client program,



❖ **Hosts** are sometimes further divided into two categories: **clients and servers.**

❖ **clients** tend to be (*desktop and mobile PCs, smartphones, and so on,*)

❖ **servers** tend to be (*more powerful machines that store and distribute Web pages, stream video, and so on*).

❖ servers often in **data centers**

For example, Google has 30–50 data centers, with many having more than one hundred thousand servers.

access networks

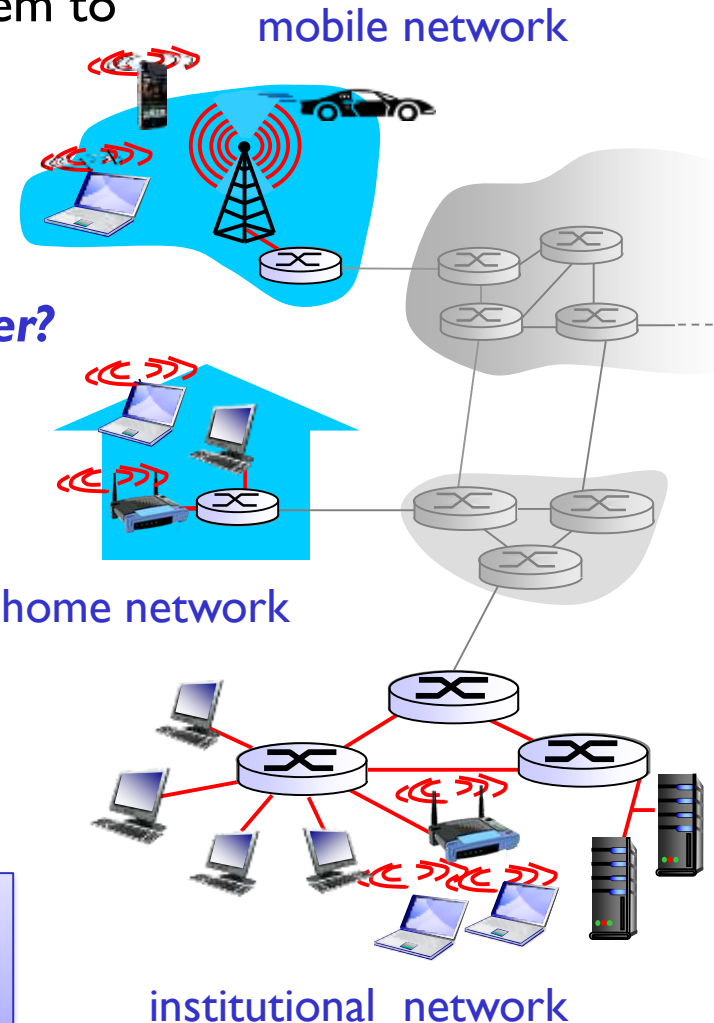
❖ the network that physically connects an end system to the first router (also known as the “edge router”)

❖ **Q: How to connect end systems to edge router?**

- home access network
- institutional access networks (school, company)
- mobile access networks

keep in mind:

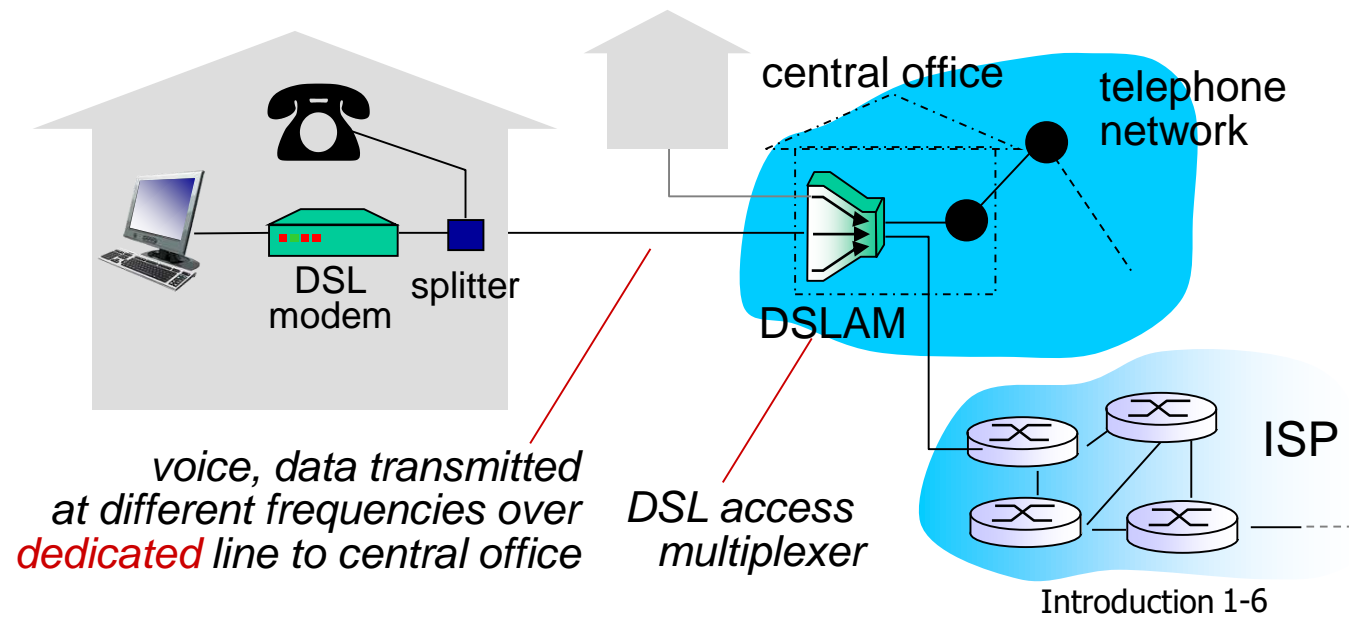
bandwidth (bits per second) of access network shared or dedicated?



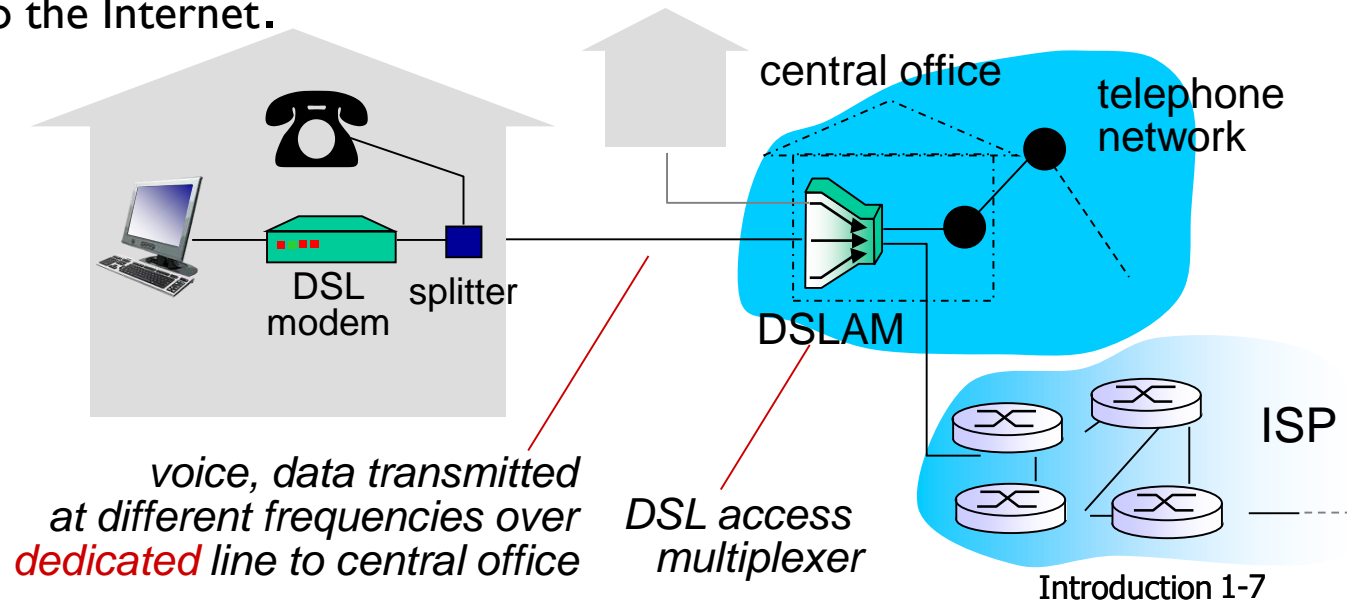
Home access nets (DSL, Cable, and Satellite)

Digital subscriber line (DSL)

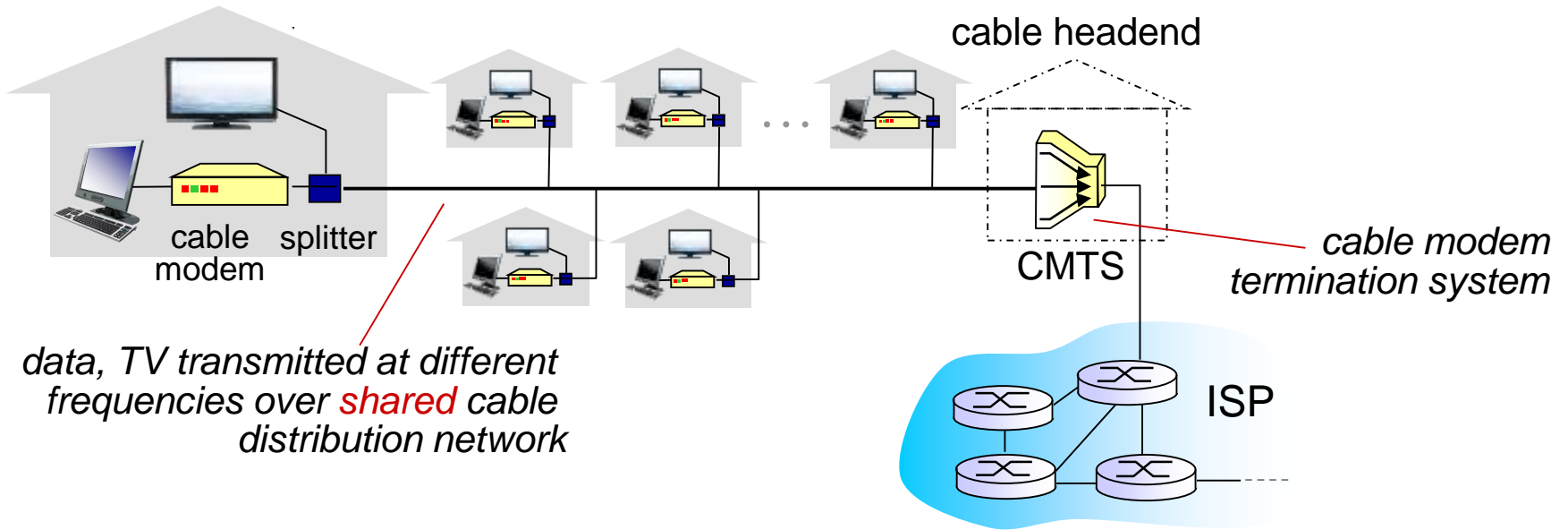
- residence typically obtains **DSL** Internet access from the same local telephone company (**telco**).
- each customer's **DSL modem** uses the existing telephone line to exchange data with a digital subscriber line access multiplexer (**DSLAM**).
- (**DSLAM**) located in the telco's local central office (CO).
- The home's **DSL modem** takes digital data and translates it to high frequency tones for transmission over telephone wires to the CO;



- the analog signals from many such houses are translated back into digital format at the **DSLAM**.
- a telephone call and an Internet connection can share the DSL link at the same time.
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
 - < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)
 - < 24 Mbps downstream transmission rate (typically < 10 Mbps)
- On the customer side, a **splitter** separates the data and telephone signals arriving to the home and forwards the data signal to the DSL modem.
- On the telco side, in the CO, the **DSLAM** separates the data and phone signals and sends the data into the Internet.



Access net: cable network

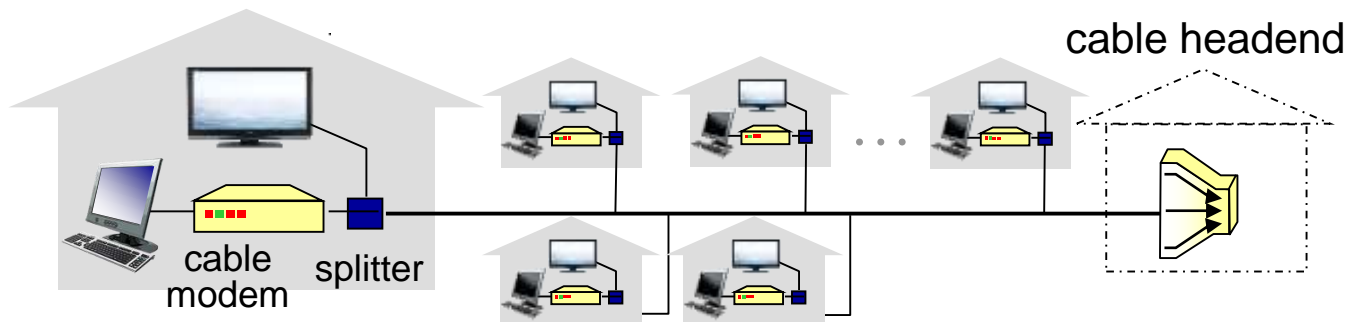


- ❖ network of cable, fiber attaches homes to ISP router
 - homes *share access network* to cable head end
 - unlike DSL, which has dedicated access to central office

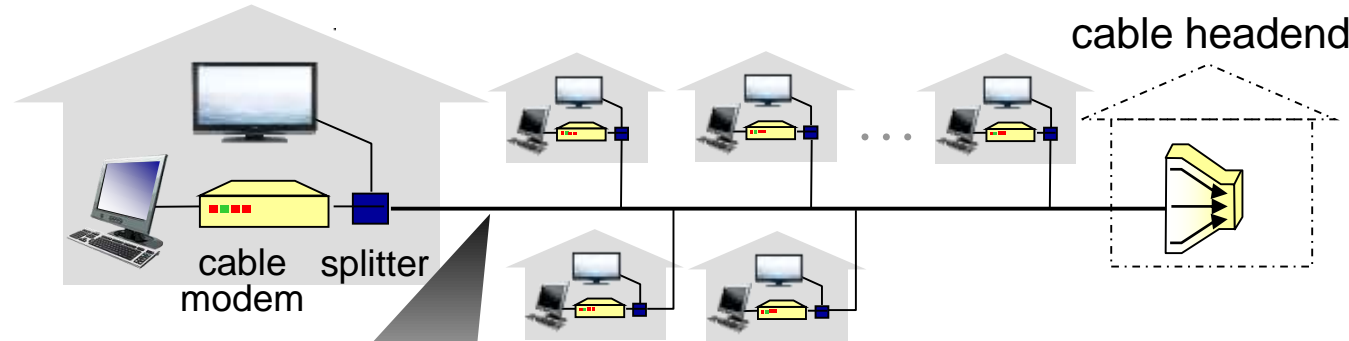
Access net: cable network

- Cable internet access requires special modems, called **cable modems**
- the cable modem is typically an external device and connects to the home PC.
- At the cable head end, the cable modem termination system (CMTS) turning the analog signal sent from the cable modems in many downstream homes back into digital format.
- Cable modems divide the HFC (**hybrid fiber coax**) network into two channels, a downstream and an upstream channel

every packet sent by the head end travels downstream on every link to every home and every packet sent by a home travels on the upstream channel to the head end.

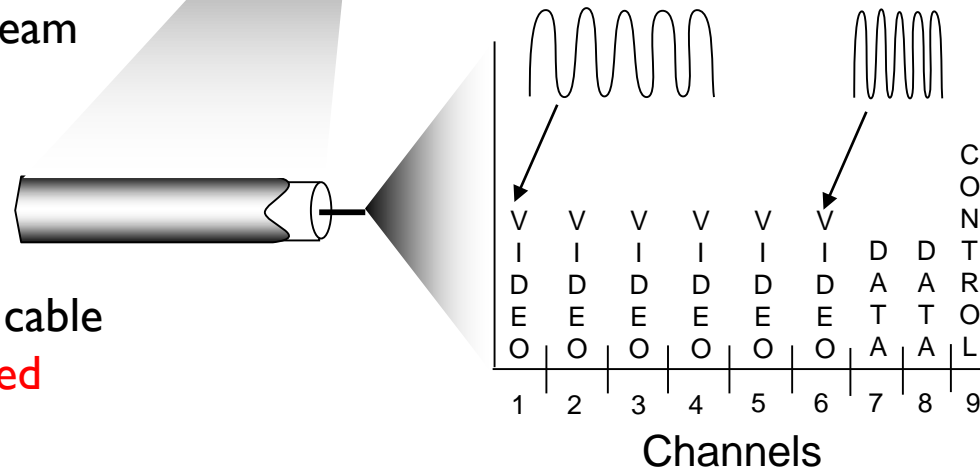


Access net: cable network



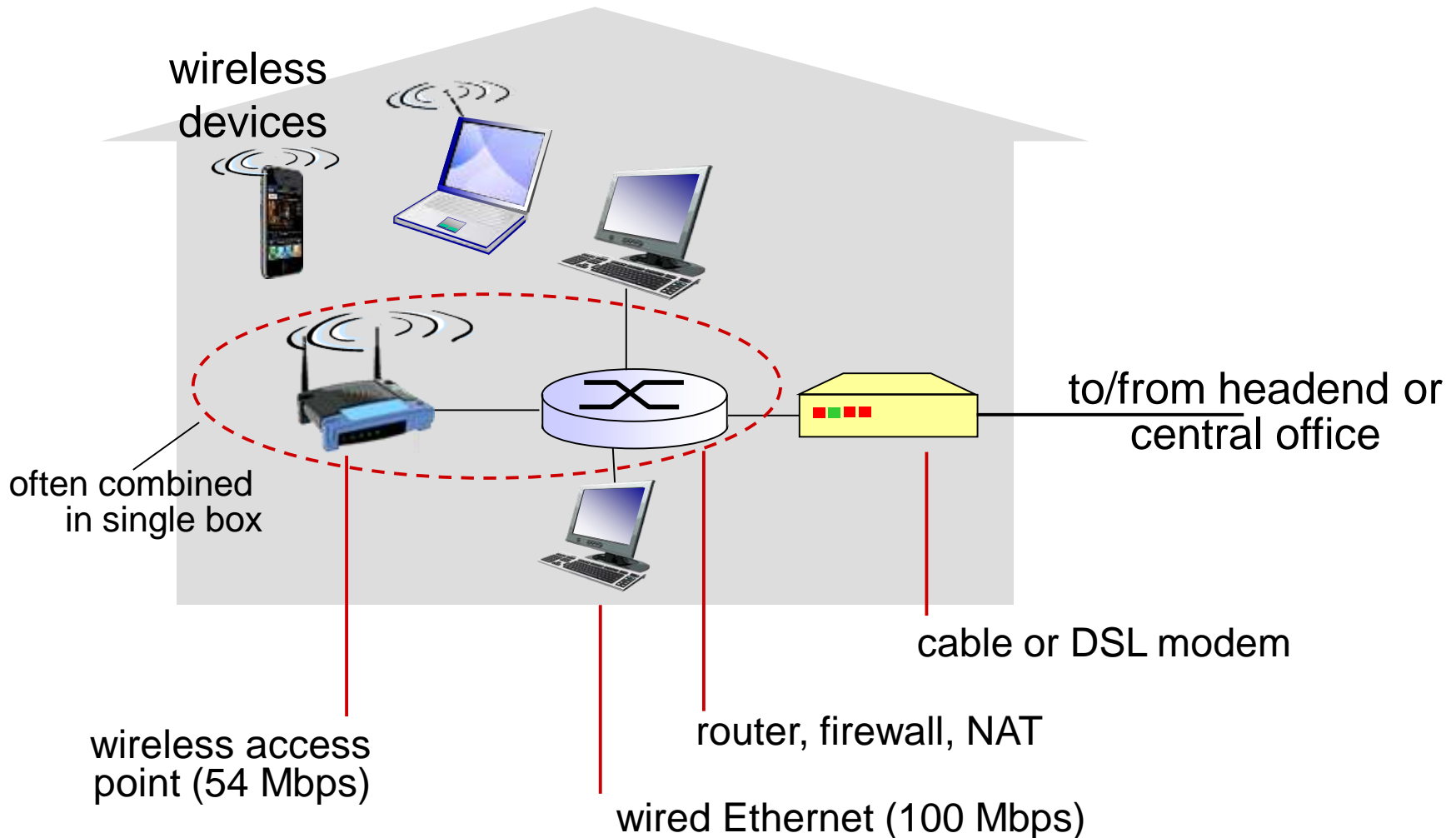
➤ asymmetric: up to 30Mbps downstream transmission rate, 2 Mbps upstream transmission rate.

➤ One important characteristic of cable Internet access is that it is a **shared** broadcast medium.

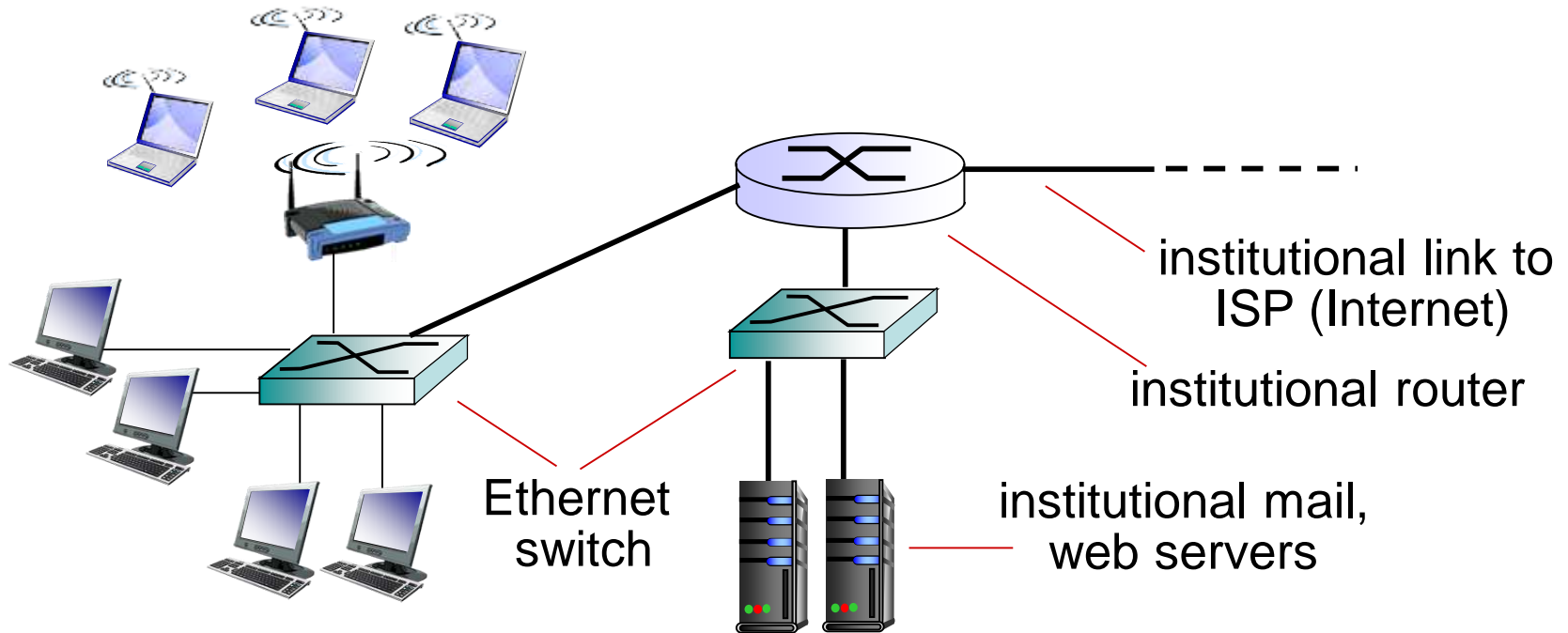


➤ **frequency division multiplexing**: different channels transmitted in different frequency bands

Access net: home network



Enterprise access networks (Ethernet)



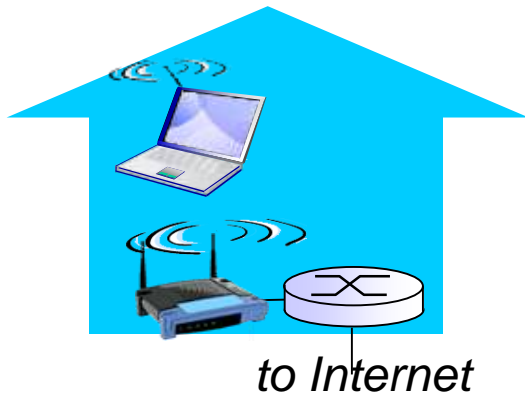
- typically used in companies, universities, etc
- The Ethernet switch or network is used to connect an end system to the edge router.
- The Ethernet switch or network, is then in turn connected into the larger Internet.
- 10 Mbps, 100Mbps, 1Gbps, 10Gbps transmission rates

Wireless access networks

- ❖ shared *wireless* access network connects end system to router
 - via base station aka “access point”

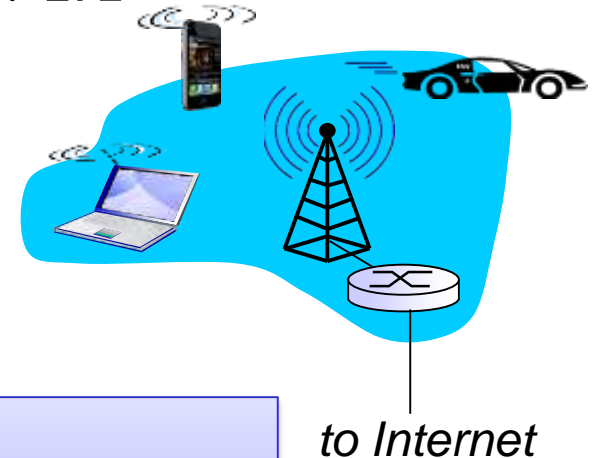
wireless LANs:

- within building (100 ft)
- 802.11b/g (WiFi): 11, 54 Mbps transmission rate



wide-area wireless access

- provided by telco (cellular) operator, 10's km
- between 1 and 10 Mbps
- 3G, 4G: LTE

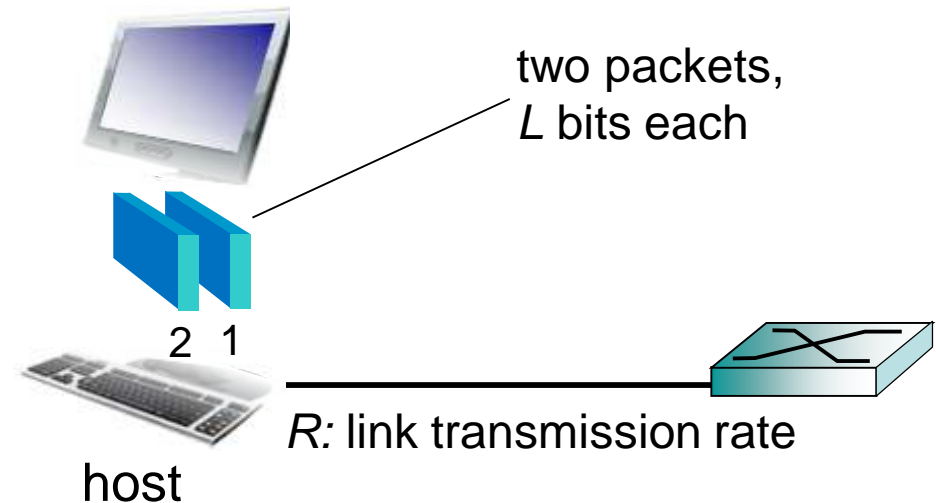


- third-generation (3G) wireless, at speeds in excess of 1 Mbps.
- higher-speed wide-area access technologies—a fourth-generation (4G)
- LTE (for “Long-Term Evolution”, and can potentially achieve rates in excess of 10 Mbps.

Host: sends *packets* of data

host sending function:

- ❖ takes application message
- ❖ breaks into smaller chunks, known as *packets*, of length L bits
- ❖ transmits packet into access network at *transmission rate R*
 - link transmission rate, aka link *capacity*, aka *link bandwidth*



$$\text{packet transmission delay} = \text{time needed to transmit } L\text{-bit packet into link} = \frac{L \text{ (bits)}}{R \text{ (bits/sec)}}$$

Transmission media

- For each transmitter-receiver pair, the bit is sent by propagating electromagnetic waves or optical pulses across a **physical medium**.
- **Physical media** fall into two categories: **guided media** and **unguided media**

With guided media, the waves are guided along a solid medium, such as a fiber-optic cable, a twisted-pair copper wire, or a coaxial cable

With unguided media, the waves propagate in the atmosphere and in outer space, such as in a wireless LAN or a digital satellite channel.

guided media

twisted pair (TP)

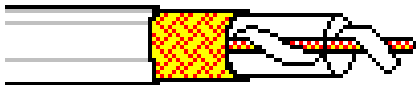
- ❖ two insulated copper wires
 - Category 5: 100 Mbps, 1 Gbps Ethernet
 - Category 6: 10Gbps
- ❖ The data rates that can be achieved depend on the thickness of the wire and the distance between transmitter and receiver.
- The least expensive
- ❖ The wires are twisted together to reduce the electrical interference .



guided media

coaxial cable:

- ❖ two concentric copper conductors
- ❖ Bidirectional
- ❖ Coaxial cable is quite common in cable television systems.
- ❖ broadband:
 - Coaxial cable can be used as a guided **shared medium**



fiber optic cable:

- ❖ glass fiber carrying light pulses, each pulse a bit
- ❖ high-speed operation:
 - high-speed point-to-point transmission (e.g., 10' s-100' s Gpbs transmission rate)
- ❖ low error rate:
 - immune to electromagnetic noise



unguided media ‘Radio Channels’

- ❖ signal carried in electromagnetic spectrum
- ❖ no physical “wire”
- ❖ can penetrate walls,
- ❖ carry a signal for long distances.

propagation environment effects:

- ❖ **reflection** (multipath fading (due to signal reflection off of interfering objects))
- ❖ **obstruction by objects** (which decrease the signal strength as the signal travels over a distance and around/through obstructing objects),
- ❖ **Interference:** (due to other transmissions and electromagnetic signals).

radio link types:

- ❖ **microwave**
 - up to 45 Mbps channels
- ❖ **LAN** (e.g., WiFi)
 - 11 Mbps, 54 Mbps
- ❖ **wide-area** (e.g., cellular)
 - 3G cellular: ~ few Mbps
- ❖ **Satellite**
 - satellite links, which can operate at speeds of hundreds of Mbps,.
 - The huge distance from **ground station** through **satellite** back to ground station introduces a substantial signal propagation delay of 280 milliseconds.

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- packet switching, circuit switching, network structure

1.4 delay, loss, throughput in networks

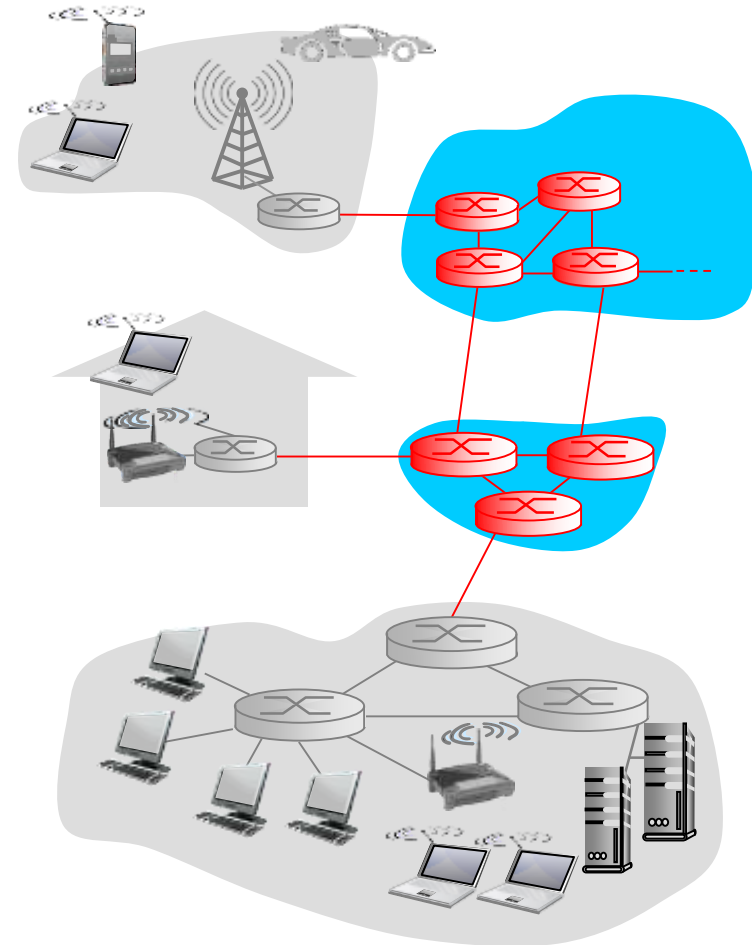
1.5 protocol layers, service models

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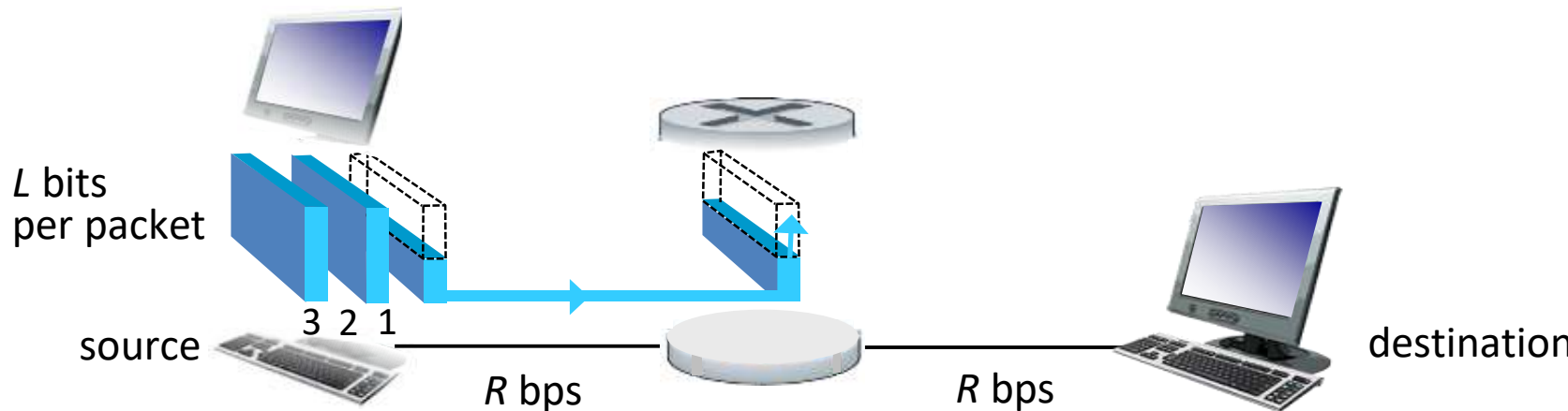
The network core

- ❖ the mesh of packet switches and links that interconnects the Internet's end systems.
- ❖ There are two fundamental approaches to moving data through a network of links and switches : (**packet switching and circuit switching .**)
- ❖ **packet-switching:**
- ❖ To send a message from a source end system to a destination end system, the source breaks long messages into smaller chunks of data known as **packets**.
- ❖ forward packets from one router to the next, across links on path from source to destination.
- ❖ each packet transmitted at full link capacity



Packet-switching: store-and-forward

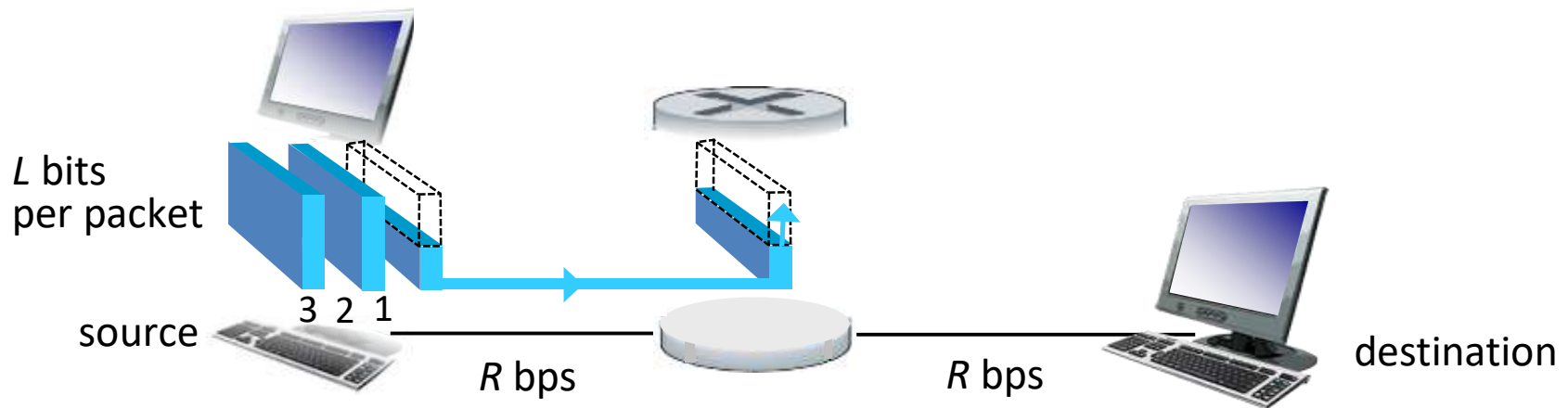
store and forward: entire packet must arrive at router before it can be transmitted on next link.



Ex:

- the router has the rather simple task of transferring a packet from one (input) link to the only other attached link.
- the **source** has three **packets**, each consisting of L bits, to send to the **destination**. the source has transmitted some of packet 1, and the front of packet 1 has already arrived at the **router**.
- the router employs **store-and-forwarding**, at this instant of time, the router cannot transmit the bits it has received, it “**store**” the packet’s bits.
- Only after the router has received *all of the packet’s bits*. can it begin to transmit “**forward**” the packet onto the outbound link.

Packet-switching: store-and-forward



- ❖ takes L/R seconds to transmit (push out) L -bit packet into link at R bps (*at the router*)
- ❖ end-end delay = $2L/R$

one-hop numerical example:

- $L = 7.5$ Mbits
- $R = 1.5$ Mbps
- one-hop transmission delay = 5 sec

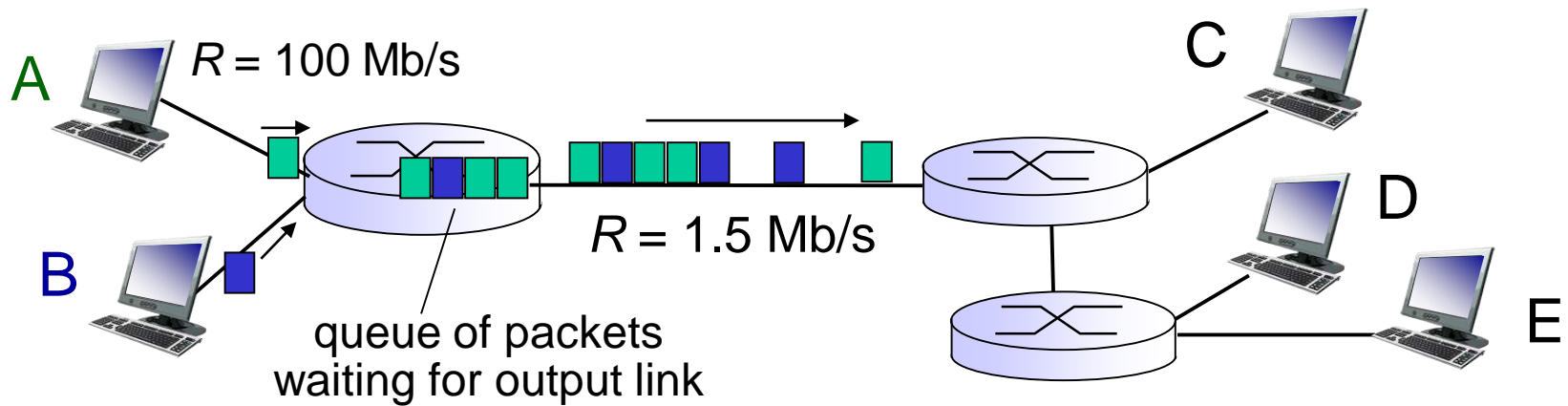
Keep In mind

end-end delay = $N L/R$

N links

there are $N-1$ routers between source and destination)

Packet Switching: queueing delay, loss



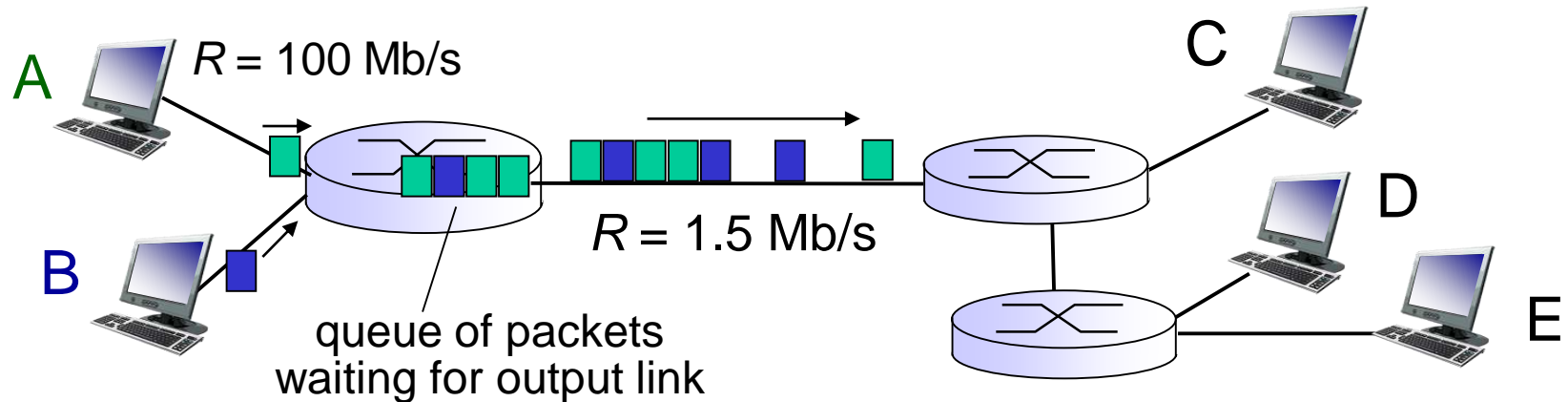
- The packet switch has an **output buffer (also called an output queue)**, which stores packets that the router is about to send into that link.
- If an arriving packet needs to be transmitted onto a link but finds the link busy with the transmission of another packet, the arriving packet must wait in the output buffer.

Packet Switching: queueing delay, loss

- in addition to the store-and-forward delays, packets suffer output buffer **queueing delays**.
- an arriving packet may find that the buffer is completely full with other packets waiting for transmission. In this case, **packet loss** will occur

queueing delays are variable and depend on the level of congestion in the network.

- In this **ex**, all packets have the same width and hence the same length
- Suppose Hosts A and B are sending packets to Host E.
- Hosts A and B first send their packets along **100 Mbps** Ethernet links to the first router.
- The router then directs these packets to the **1.5 Mbps** link.



- congestion will occur at the router as packets queue in the link's output buffer before being transmitted onto the link.

queuing and loss:

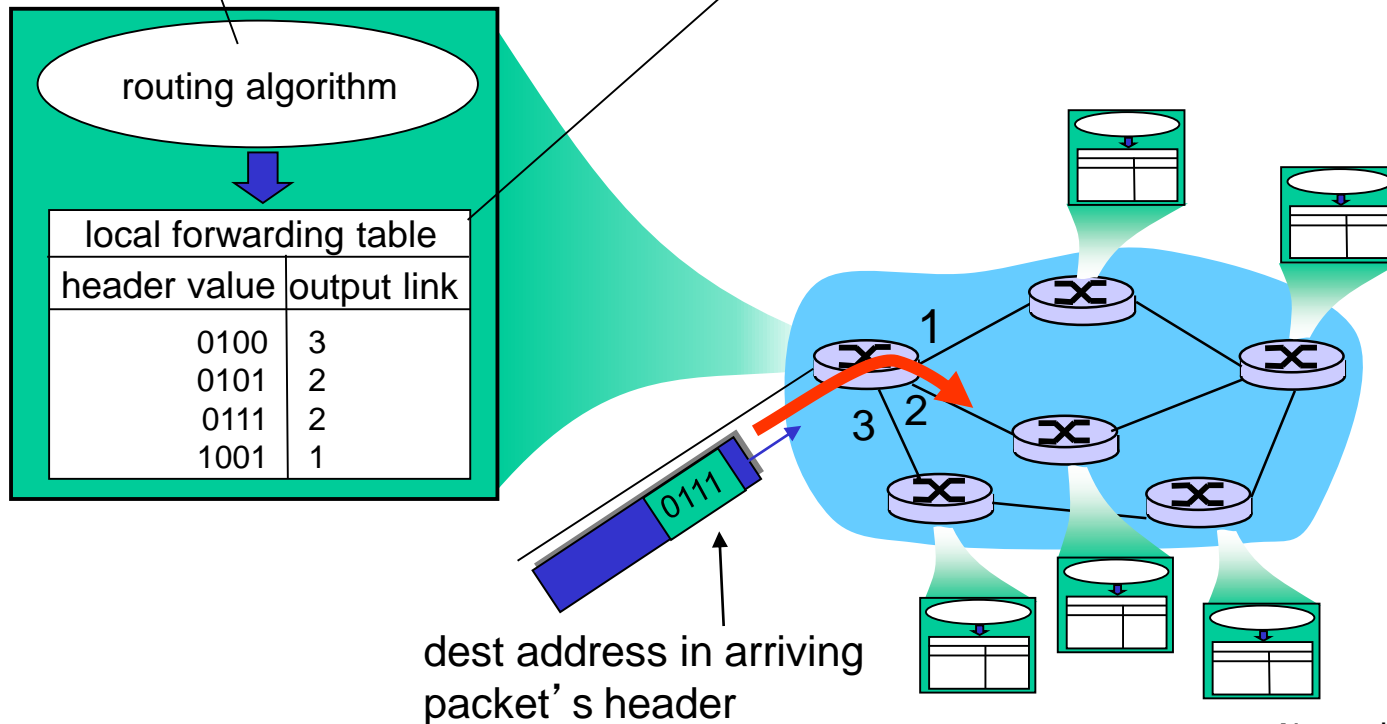
- ❖ If arrival rate (in bits) to link exceeds transmission rate of link for a period of time:
 - packets will queue, wait to be transmitted on link
 - packets can be dropped (lost) if memory (buffer) fills up

Two key network-core functions

routing: determines source-destination route taken by packets

- *routing algorithms*

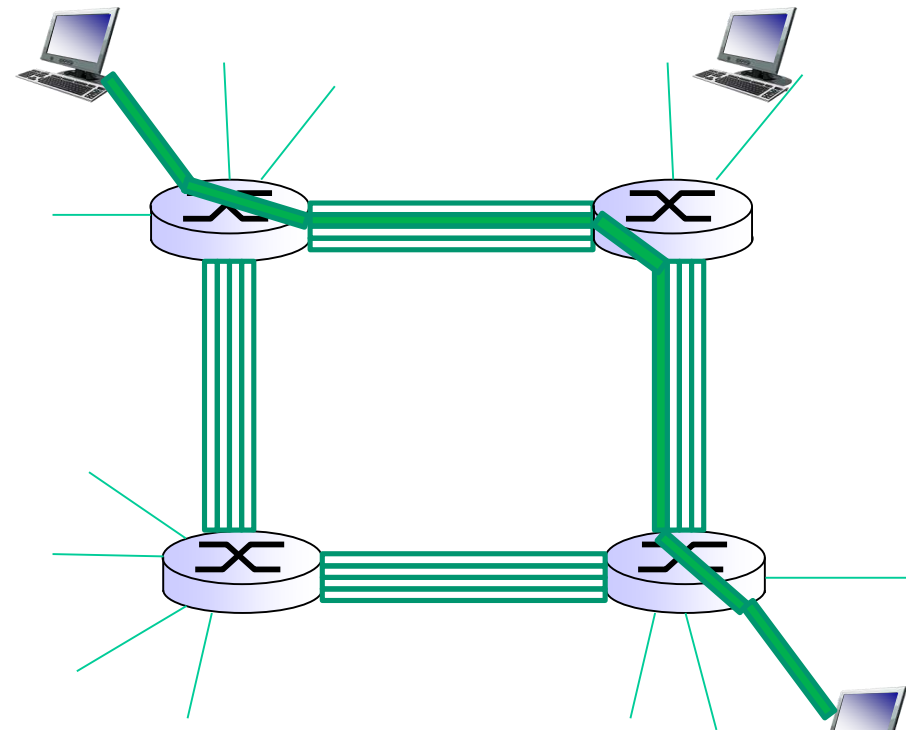
forwarding: move packets from router's input to appropriate router output



Alternative core: circuit switching

end-end resources allocated to, reserved for “call” between source & dest:

- ❖ In diagram, each link has four circuits.
 - call gets 2nd circuit in top link and 1st circuit in right link.
- ❖ dedicated resources: no sharing
 - circuit-like (guaranteed) performance
- ❖ circuit segment idle if not used by call (*no sharing*)
- ❖ Commonly used in traditional telephone networks

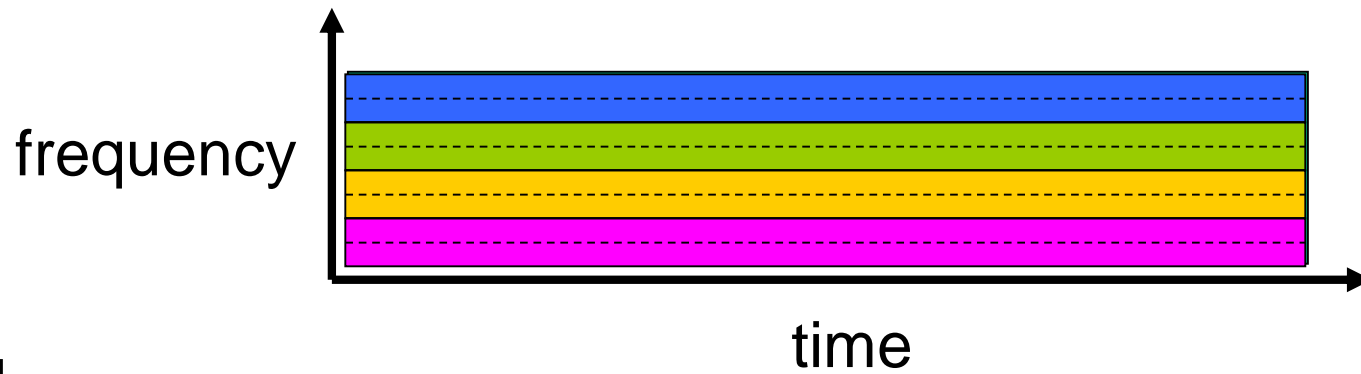


Circuit switching: FDM versus TDM

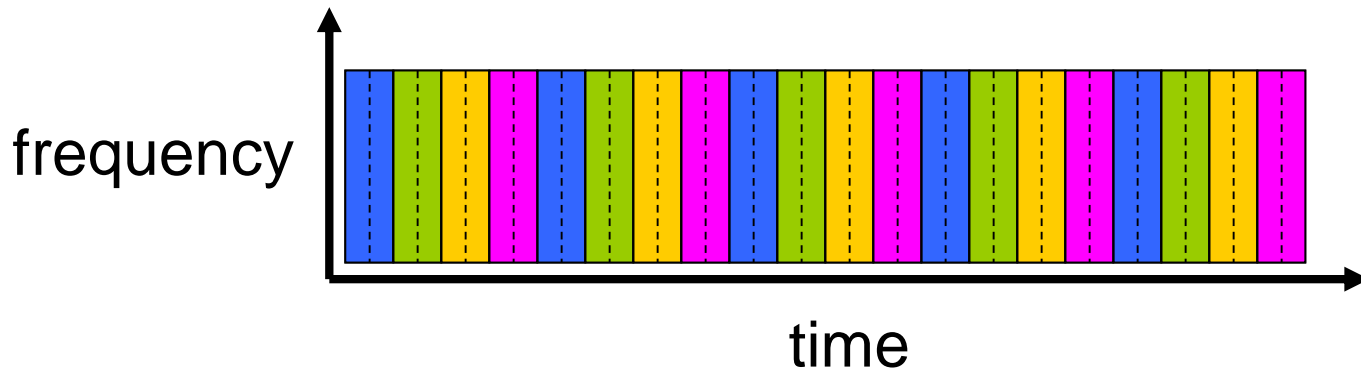
FDM

Example:

4 users



TDM

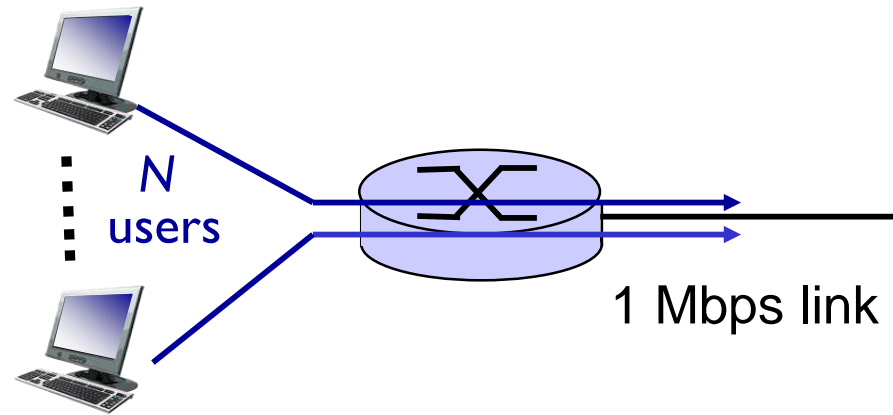


Packet switching versus circuit switching

packet switching allows more users to use network!

example:

- 1 Mb/s link
- each user:
 - 100 kb/s when “active”
 - active 10% of time



❖ *circuit-switching:*

- 10 users

❖ *packet switching:*

- with 35 users, probability > 10 active at same time is less than $.0004$ *

* Check out the online interactive exercises for more examples

Packet switching versus circuit switching

is packet switching a “slam dunk winner?”

- ❖ great for bursty data
 - resource sharing
 - simpler, no call setup
- ❖ **excessive congestion possible:** packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- ❖ **Q: How to provide circuit-like behavior?**
 - bandwidth guarantees needed for audio/video apps
 - still an unsolved problem (chapter 7)

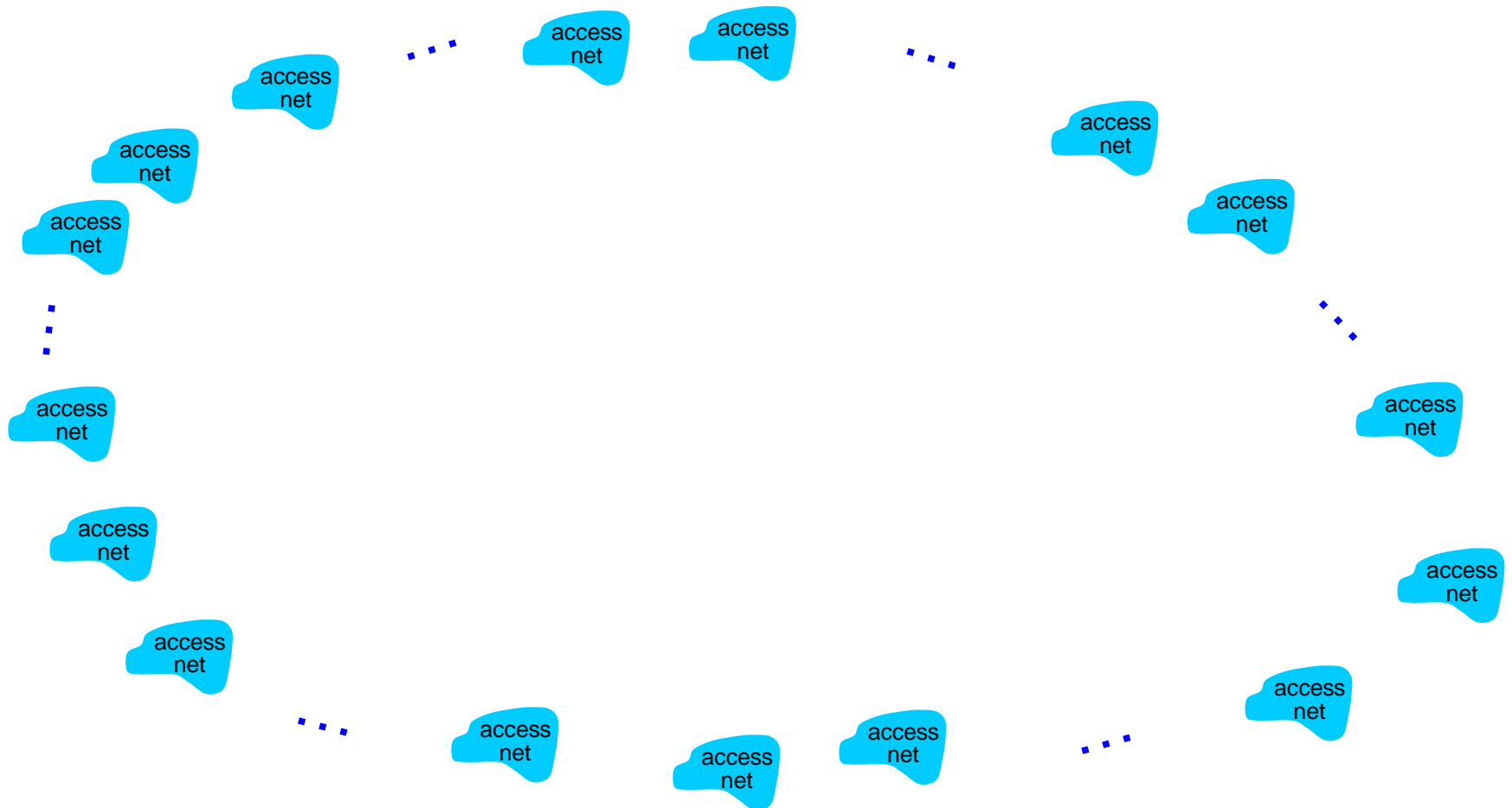
Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?

Internet structure: network of networks

- ❖ End systems connect to Internet via **access ISPs** (Internet Service Providers)
 - Residential, company and university ISPs
- ❖ Access ISPs in turn must be interconnected.
 - ❖ So that any two hosts can send packets to each other
- ❖ Resulting network of networks is very complex
 - ❖ Evolution was driven by **economics** and **national policies**
- ❖ Let's take a stepwise approach to describe current Internet structure

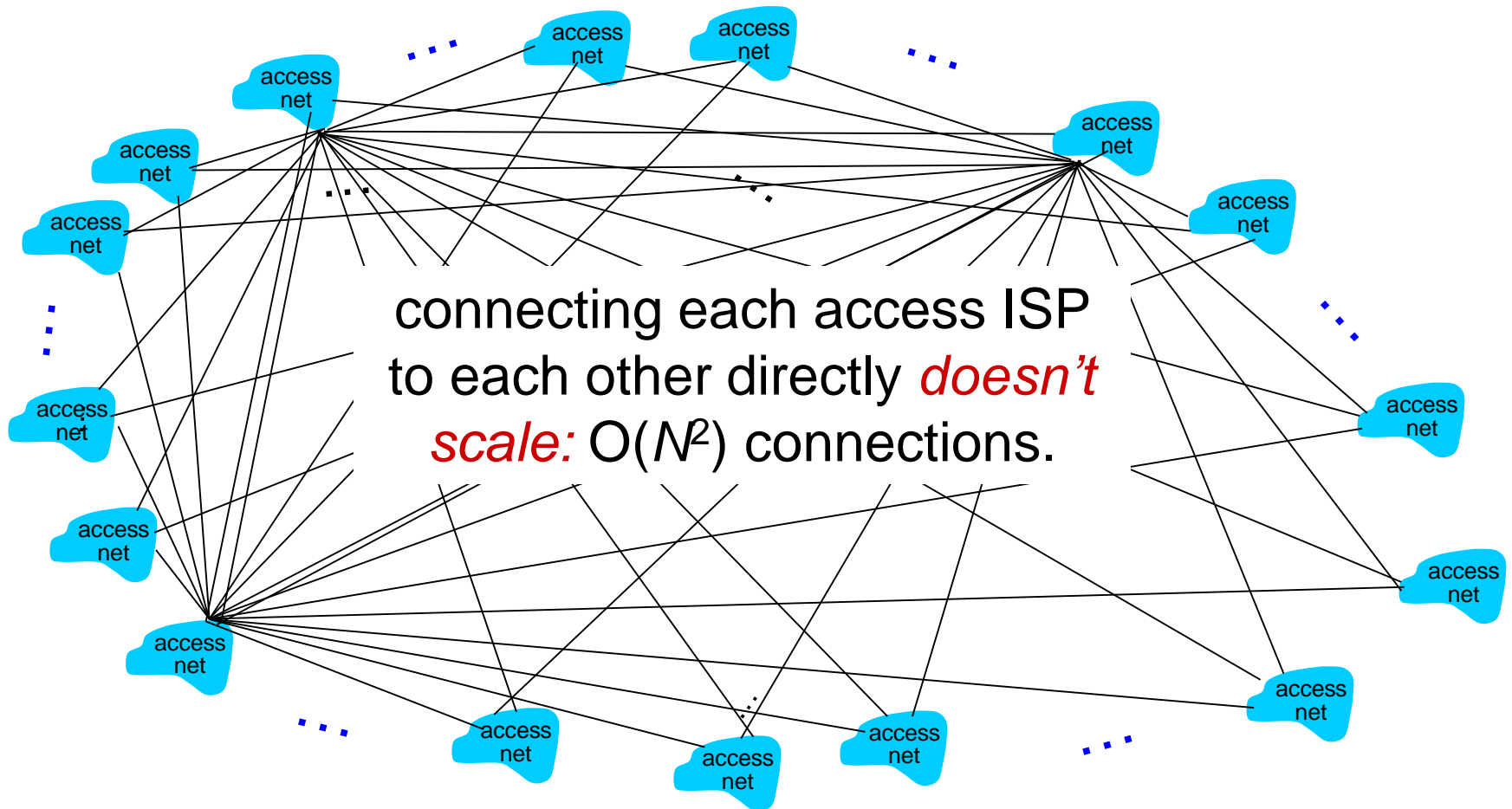
Internet structure: network of networks

Question: given *millions* of access ISPs, how to connect them together?



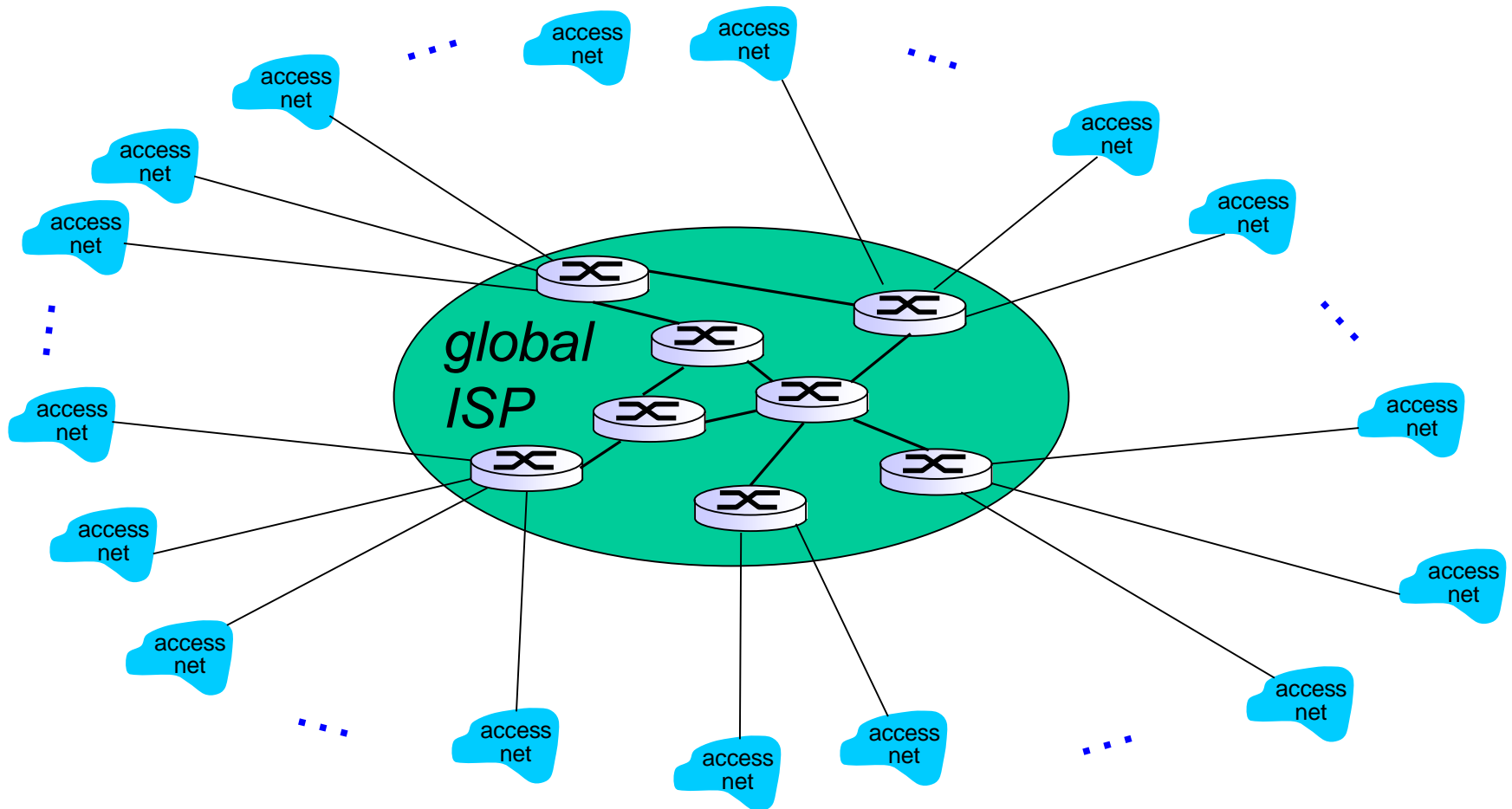
Internet structure: network of networks

Option: connect each access ISP to every other access ISP?



Internet structure: network of networks

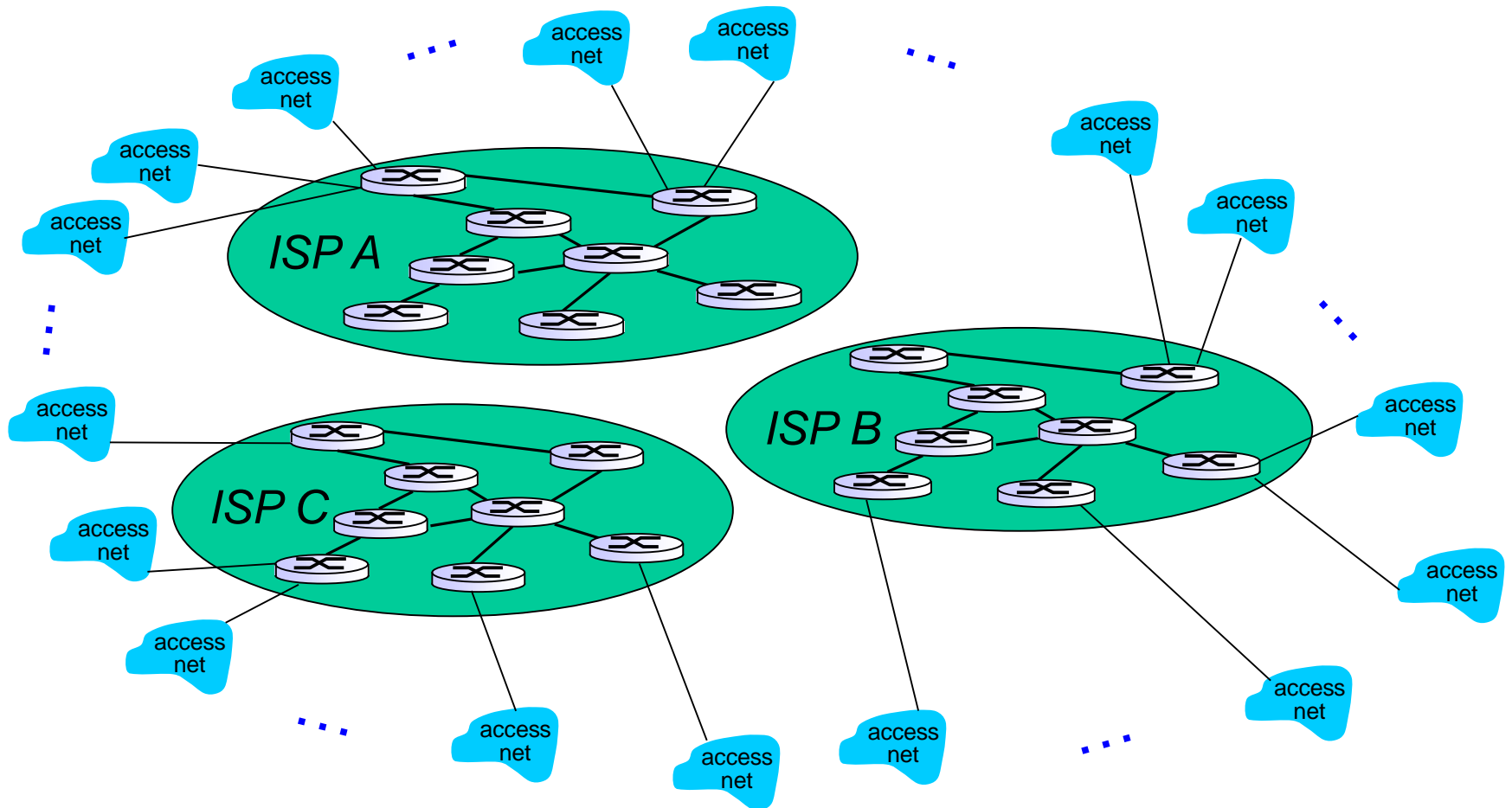
Option: connect each access ISP to a global transit ISP? Customer and provider ISPs have economic agreement.



Internet structure: network of networks

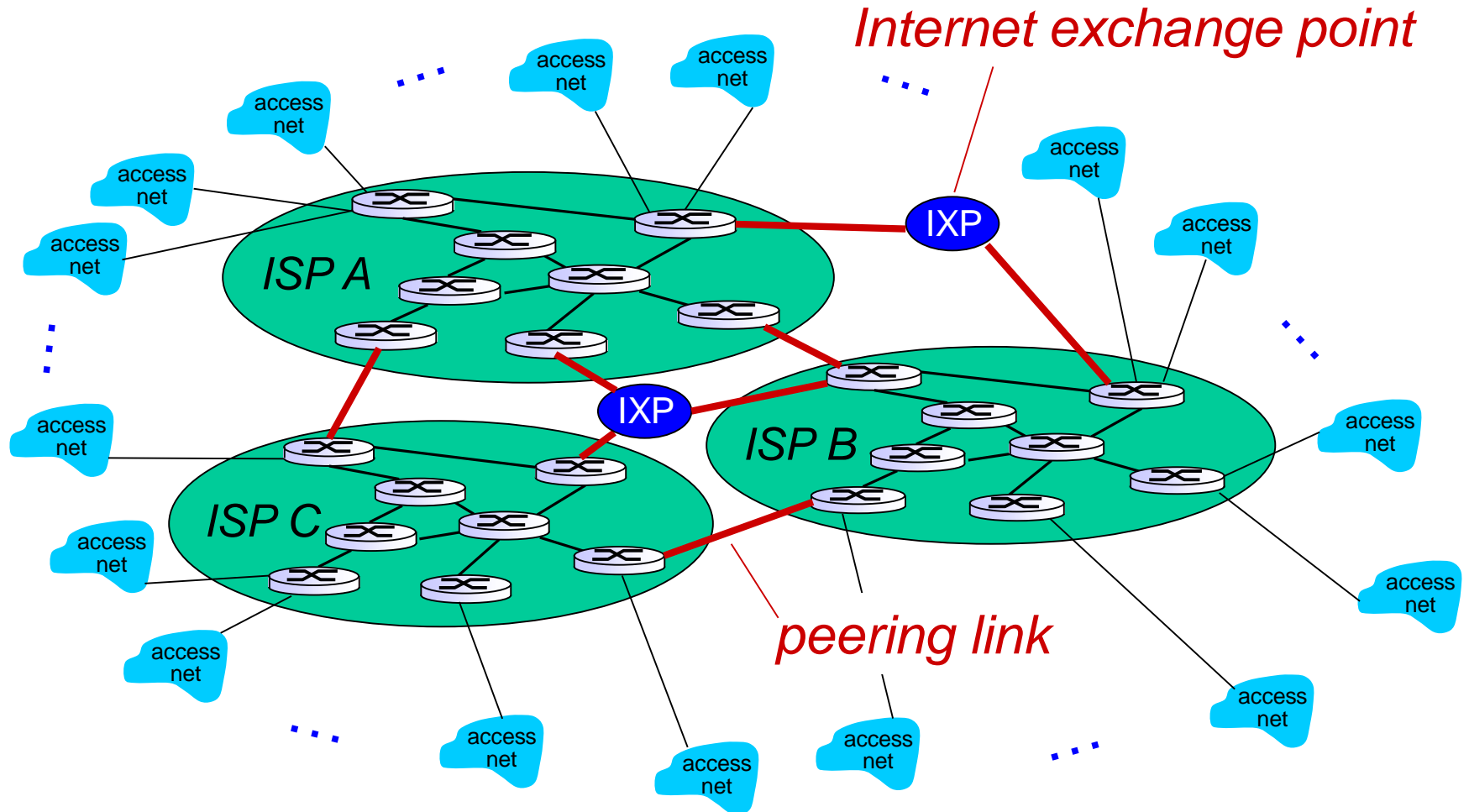
But if one global ISP is viable business, there will be competitors

....



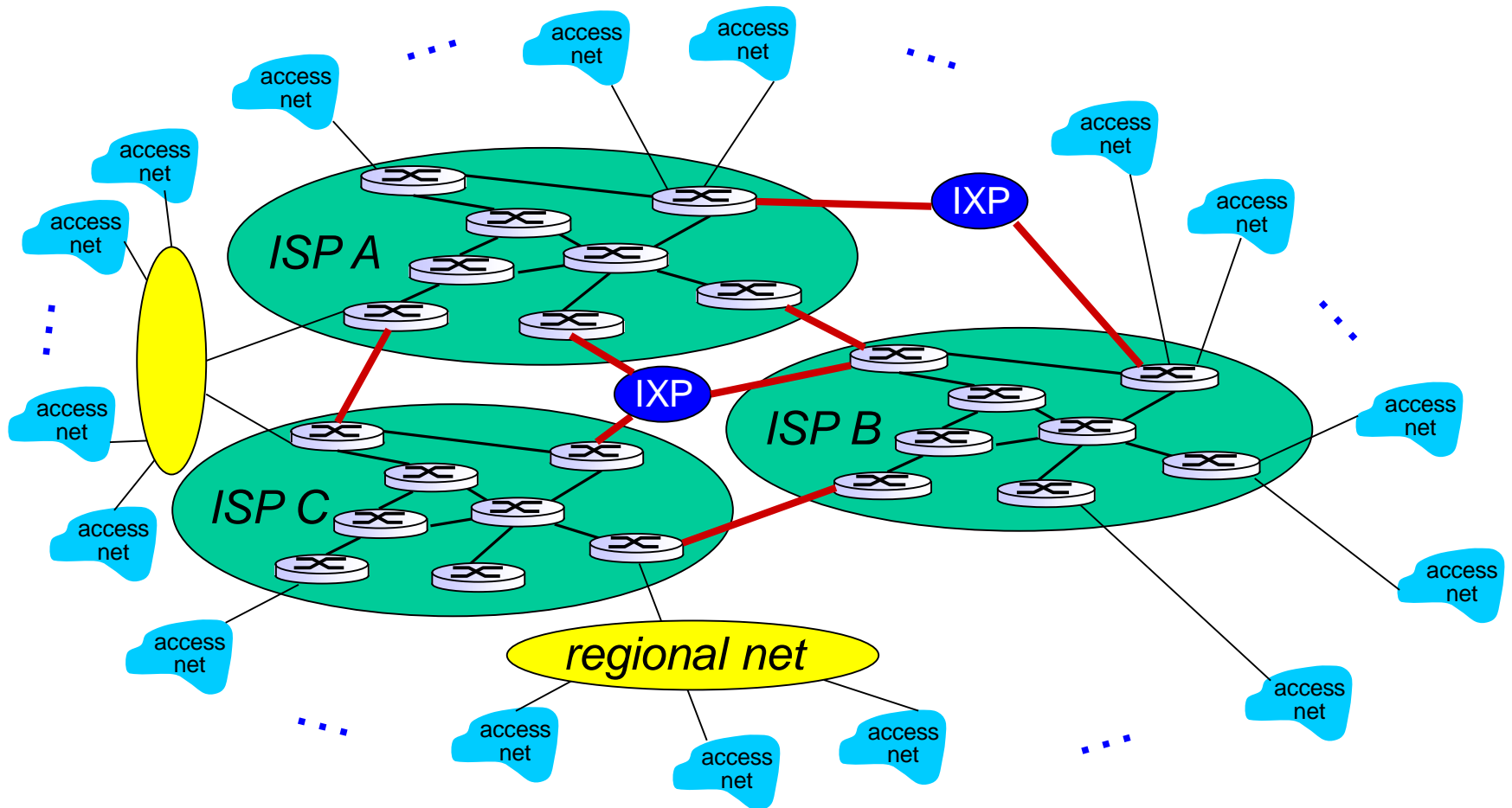
Internet structure: network of networks

But if one global ISP is viable business, there will be competitors
.... which must be interconnected



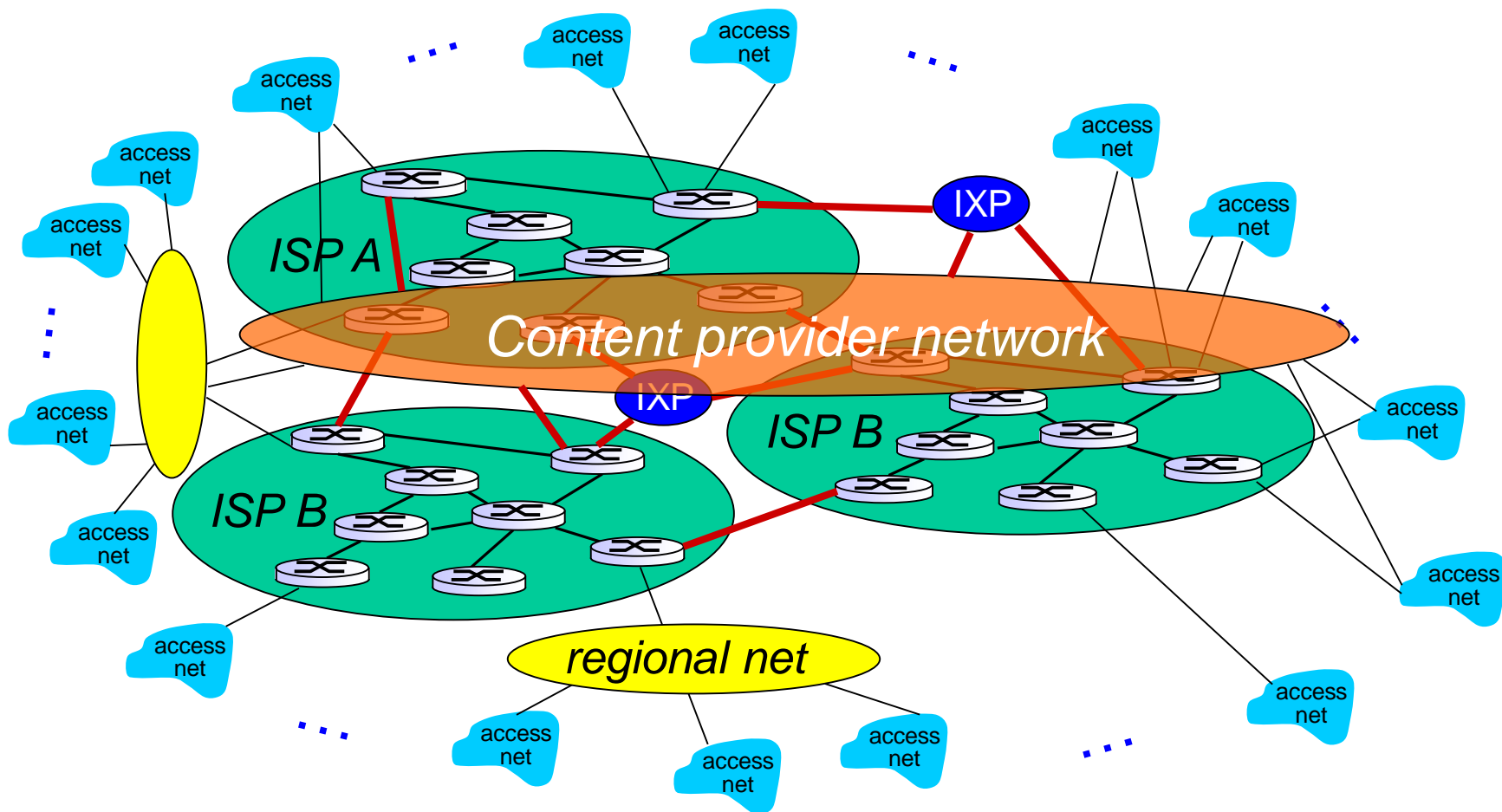
Internet structure: network of networks

... and regional networks may arise to connect access nets to ISPS

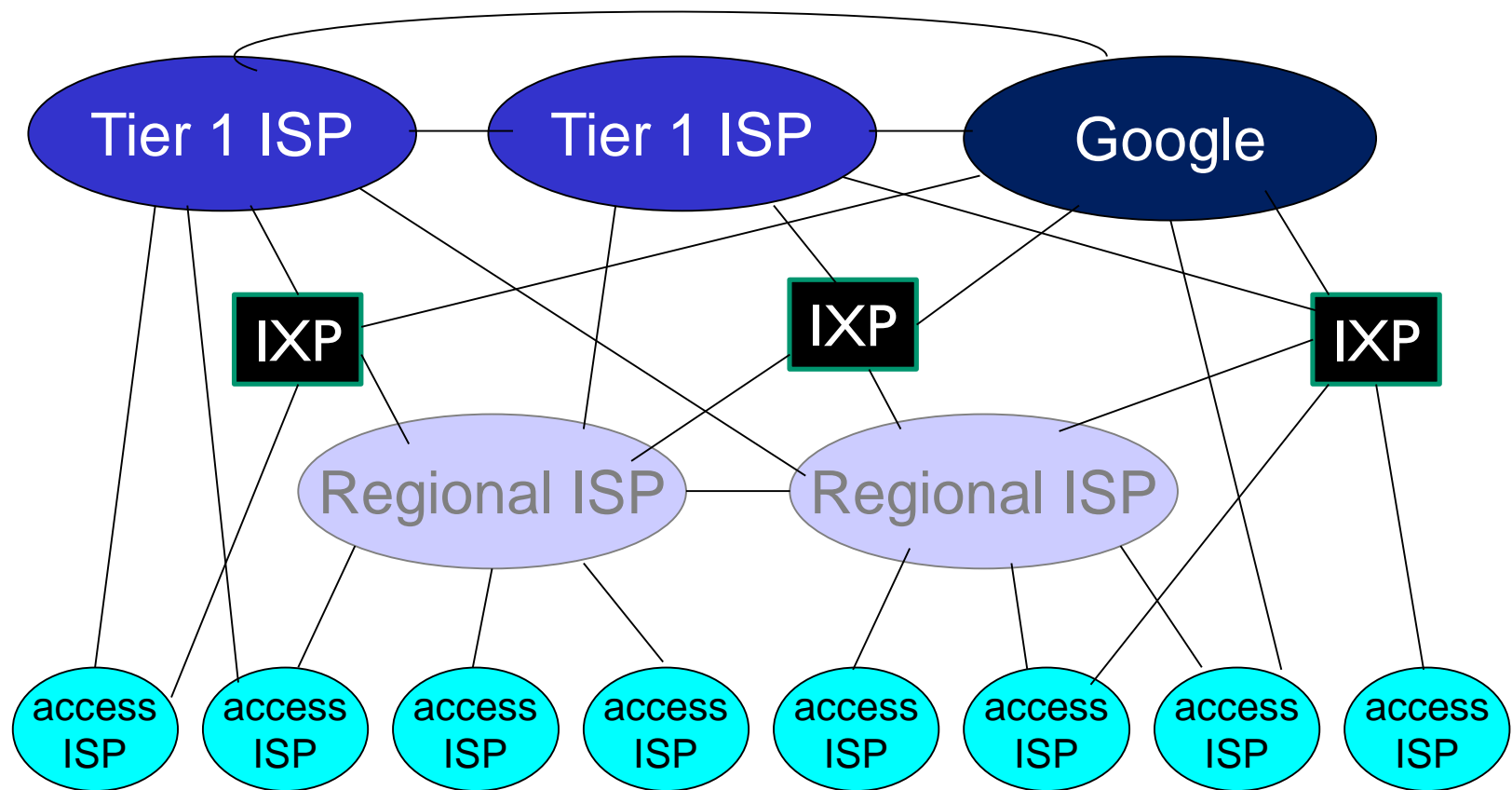


Internet structure: network of networks

... and content provider networks (e.g., Google, Microsoft, Akamai) may run their own network, to bring services, content close to end users



Internet structure: network of networks



- ❖ at center: small # of well-connected large networks
 - “tier-1” commercial ISPs (e.g., Level 3, Sprint, AT&T, NTT), national & international coverage
 - content provider network (e.g., Google): private network that connects its data centers to Internet, often bypassing tier-1, regional ISPs

Tier-1 ISP: e.g., Sprint

