



## *Design and Analysis Algorithms*

### تصميم وتحليل خوارزميات

**ITGS301**

المحاضرة العاشرة: Lecture 10



## Graph Algorithms

### What is a Graph?

A Graph is an abstract data structure that represents a collection of items with pairwise relationship between these items.

It consists of a set of vertices (nodes) connected by a set of edges (links), and is denoted by  $G = (V, E)$ , where  $V$  is set of vertices and  $E$  is set of edges.



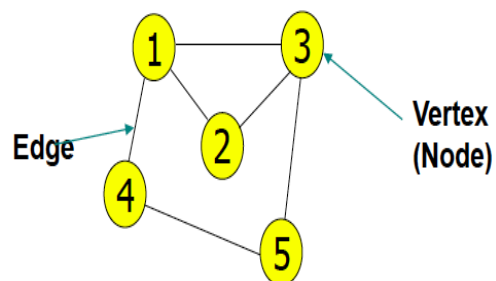
- $V(G)$  is a set of vertices or nodes which can represent an object that needs to be “connected”.
- $V$  represents the number of vertices ( nodes) in the graph
- $E(G)$  is a set of edges. An edge is a distinct pair of vertices. An edge indicates a valid/existing connection between two vertices.
- $E$  represents the number of edges in the graph



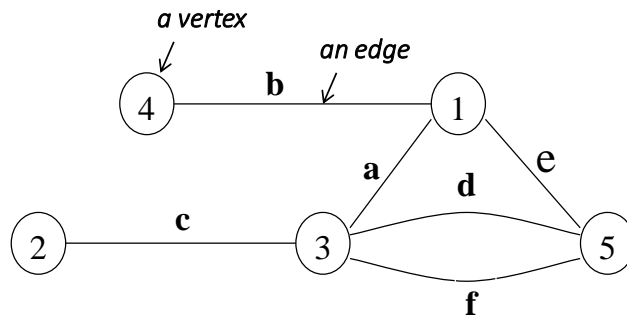
$$G = (V, E)$$

$V$  = set of vertices     $|V| = n$

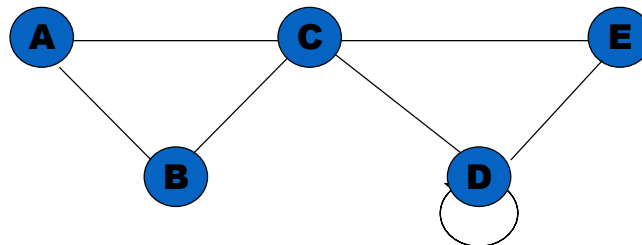
$E$  = set of edges     $|E| = m$



## An example of a graph



## Another Example



$V = 5$

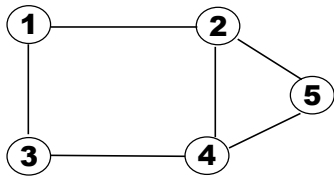
$V(G) = \{A, B, C, D, E\}$

$E(G) = \{(AC), (AB), (BC), (CD), (CE), (DE), (DD)\}$



## Graph representation

1. **Adjacency Matrix:** represent a graph as  $n \times n$  Matrix A

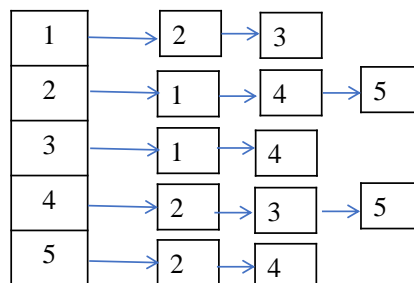
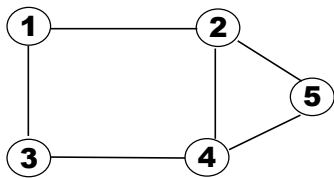


|   | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|
| 1 | 0 | 1 | 1 | 0 | 0 |
| 2 | 1 | 0 | 0 | 1 | 1 |
| 3 | 1 | 0 | 0 | 1 | 0 |
| 4 | 0 | 1 | 1 | 0 | 1 |
| 5 | 0 | 1 | 0 | 1 | 0 |

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## Graph representation

2. **Adjacent List:** for each vertex  $v \in V$ , store a list of vertices adjacent to  $v$

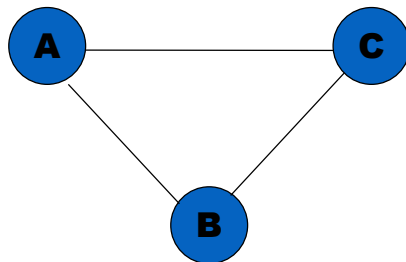


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## Graph Terminology

- Adjacent Vertices

if two vertices are joined by an edge they are said to be **adjacent**



Adjacent Vertices:

A & C

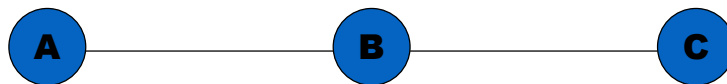
A & B

B & C

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- Degree

- the degree of a vertex  $x$  is the number of vertices adjacent to it (or the number of edges incident to it)
- represented as  **$deg(x)$**



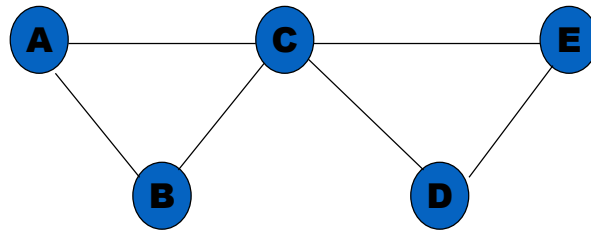
$$deg(A) = 1$$

$$deg(B) = 2$$

$$deg(C) = 1$$

## • Path

- a path is sequence of vertices in which each vertex is adjacent to the next one.

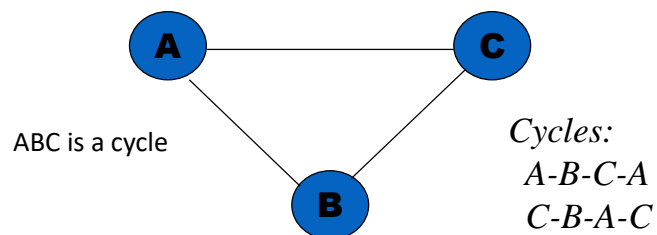


*Path from A to D: A-B-C-D*

*Path from B to E: B-A-C-D-E*

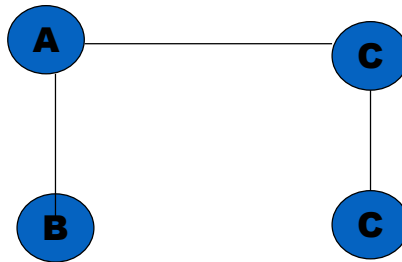
## • Cyclic

- cycle is a path consisting of *at least three vertices* that started and ends with the same vertex.
- So the graph is a cycle if there is subgraph is cycle.



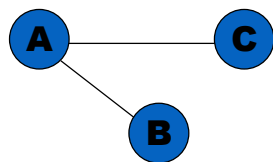
- **Acyclic**

A graph is acyclic if no subgraph is a cycle >

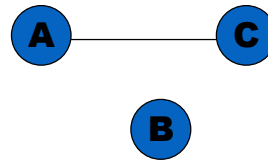


- **Connected**

– a graph  $G$  is connected if there is at least one path from every vertex to every other vertex in the graph .



Connected Graph

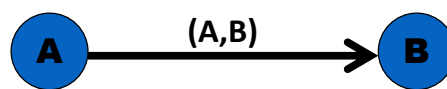


Disconnected Graph

## Types of Graph

- Directed Graph or Digraph

– the connecting lines are usually represented with an arrow

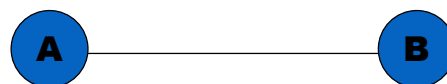


*Note:*  $(A,B) \neq (B,A)$

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- Undirected Graph

– the order of the vertices in the pair of vertices in the set of edges does not matter



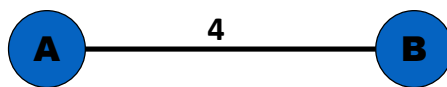
$(A,B) = (B,A)$

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## • Weighed Graph

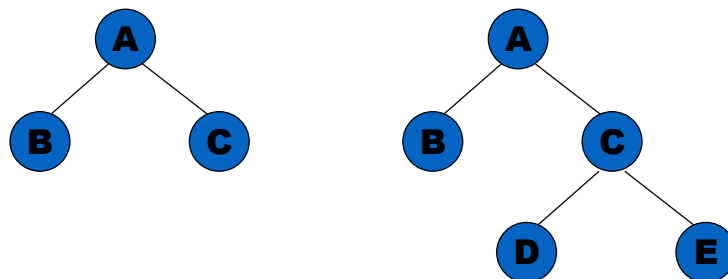
- each edge has an associated weight which could indicate cost, distance, time, etc. between two adjacent vertices



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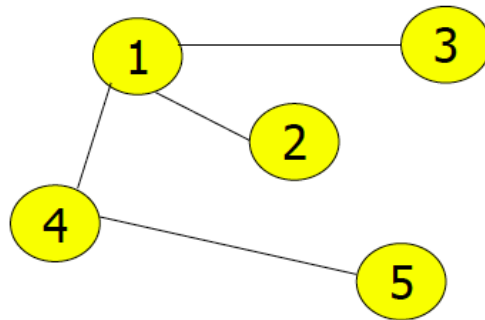
## • Tree

- Definition: A tree is a connected undirected graph with no simple circuits.
- a connected graph with no cycles



*Examples of a Tree*

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Tree

## Subgraph

- Suppose that  $V(G)$  and  $E(G)$  denote the vertex and edge sets of a graph  $G$ . If  $H$  is graph with the properties:

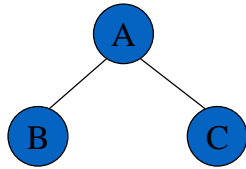
$$V(H) \subseteq V(G)$$

$$E(H) \subseteq E(G)$$

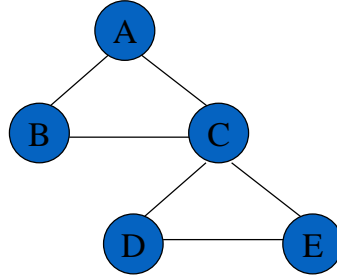
Every edge of  $E(H)$  has both its incident vertices in  $V(H)$  then  $H$  is called a **subgraph of  $G$**

# Subgraph

Graph A



Graph B

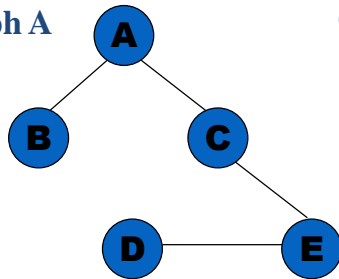


*A is a subgraph of B*

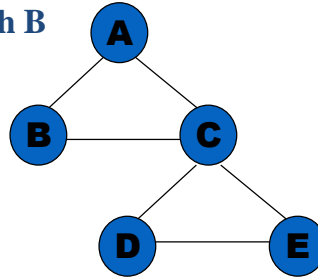
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- If  $V(H)=V(G)$  then  $H$  is called a **spanning subgraph** of  $G$

Graph A

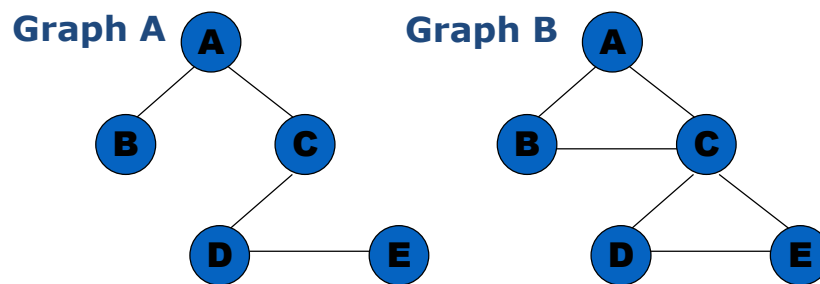


Graph B



*A is a spanning subgraph of B*

- If  $H$  is a tree, then  $H$  is called a **spanning tree**



*A is a spanning tree of B*

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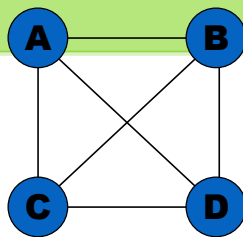
## Spanning Trees

- A spanning tree of a graph is just a sub graph that contains all the vertices and is a tree.
- A graph may have many spanning trees.
- for each graph  $G$  with  $n$  vertices , any spanning tree must has  $n-1$  edges, and *no cycle* on it.

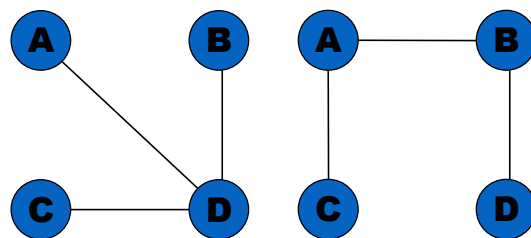
## Spanning Tree properties:

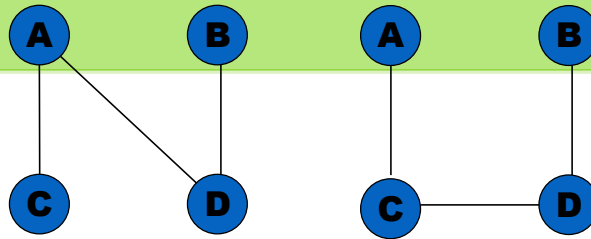
On a connected graph  $G=(V, E)$ , a spanning tree *must be:*

- a connected subgraph (contains all vertices of  $G$ )
- no cycle.
- is a tree ( $|E| = |V| - 1$ )

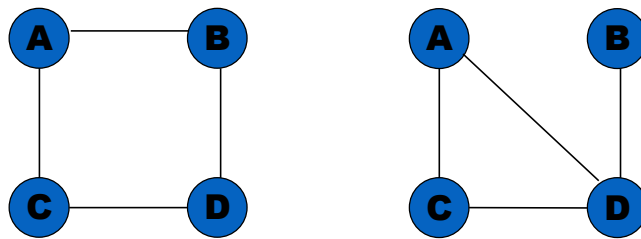


A connected undirected *graph G*





Four of the *spanning trees* of the graph



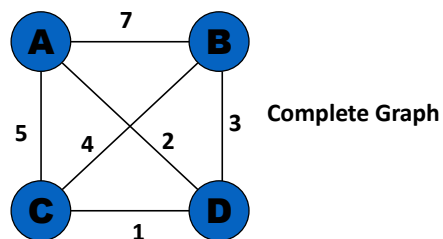
The two *are not* spanning tree

## Minimum Spanning Tree

- The Minimum Spanning Tree for a given graph is the Spanning Tree of minimum cost for that graph .
- a minimum spanning tree (MST) is a spanning tree of minimum weight

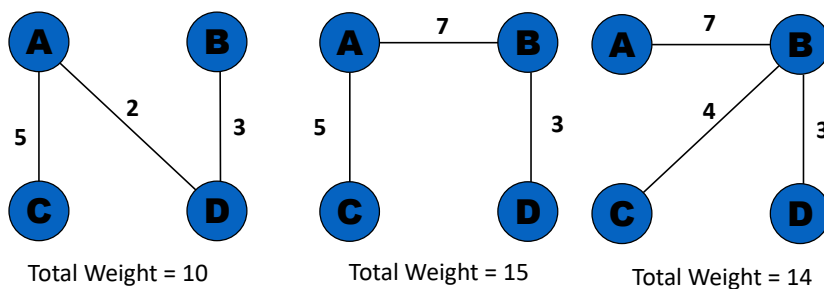
**Note** : we need to have spanning tree that connected to all its vertices but has less weights.

**Example:**



**Note:** Number of spanning tree of complete graph =  $n^{n-2}$

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All These subgraphs are spanning tree, all its vertices are connected and there is no cycle. **But, they are not minimum spanning tree because the total weights are not the least total weight.**

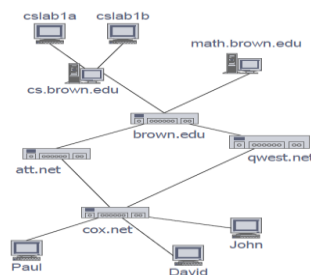
## MST Algorithms

- Minimum Spanning Tree
  - Kruskal's Algorithm
  - Primm's Algorithm

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## Applications

- ◆ **Computer networks –**
  - Local area network
  - Internet
- ◆ **Transportation networks**
  - highway network
  - flight network
- ◆ **communication networks**







*The End.* 

